April 12, 2023

Exercise 1 (Dual formulation). Let $g : \mathbb{R}^n \to \mathbb{R}^m$. Show that

1. $\mathbb{I}_{g(x)=0} = \sup_{\lambda \in \mathbb{R}^m} \lambda^\top g(x)$

2.
$$\mathbb{I}_{g(x)\leq 0} = \sup_{\lambda\in\mathbb{R}^m_+} \lambda^+ g(x)$$

3. $\mathbb{I}_{g(x)\in C} = \sup_{\lambda\in -C^{\oplus}} \lambda^{\top} g(x)$ where C is a closed convex cone, and $C^{\oplus} := \{\lambda \in \mathbb{R}^m \mid \lambda^{\top} c \geq 0, \forall c \in C\}.$

Exercise 2 (Linear Programming). Consider the following linear problem (LP)

$$(P) \quad \underset{x \ge 0}{\operatorname{Min}} \quad c^{\top} x$$
$$s.t. \quad Ax =$$

b

- 1. Show that the dual of (P) is an LP.
- 2. Show that the dual of the dual of (P) is equivalent to (P).

Exercise 3 (Quadratically Constrained Quadratic Programming). *Consider the problem*

$$(QCQP) \quad \underset{x \in \mathbb{R}^n}{\operatorname{Min}} \quad \frac{1}{2}x^{\top}P_0x + q_0^{\top}x + r_0$$
$$\frac{1}{2}x^{\top}P_ix + q_i^{\top}x + r_i \leq 0 \quad \forall i \in \mathbb{R}^n$$

where $P_0 \in S_{++}^n$ and $P_i \in S_{+}^n$.

- 1. Show by duality that, for $\mu \in \mathbb{R}^m_+$, there exists P_{μ}, q_{μ} and r_{μ} , such that $g(\mu) = -\frac{1}{2}q_{\mu}P_{\mu}^{-1}q_{\mu} + r_{\mu} \leq val(P).$
- 2. Give an easy condition under which $val(P) = \sup_{\mu \ge 0} g(\mu).$

Exercise 4 (Conic Programming). Let $K \subset \mathbb{R}^n$ be a closed convex pointed cone, and denote $x \preceq_K y$ iff $y \in x + K$. Consider the following program, with $A \in M_{m,n}$ and $b \in \mathbb{R}^m$.

$$(P) \quad \underset{x \in \mathbb{R}^n}{\min} \quad c^{\top} x$$
$$s.t. \quad Ax = b$$
$$x \preceq_K 0$$

- 1. Show that (P) is a convex optimization problem.
- 2. Denote $\mathcal{L}(x,\lambda,\mu) = c^{\top}x + \lambda^{\top}(Ax b) + \mu^{\top}x$. Show that $\operatorname{val}(P) = \operatorname{Min}_{x \in \mathbb{R}^n} \sup_{\lambda \in \mathbb{R}^m, \mu \in K^{\oplus}} \mathcal{L}(x,\lambda,\mu)$.
- 3. Give a dual problem to (P).

Exercise 5 (Duality gap). Consider the following problem

$$\begin{array}{ll} & \mathop{\rm Min}_{x\in\mathbb{R},y\in\mathbb{R}^+_*} & e^{-x} \\ & s.t. & x^2/y \leq 0 \end{array}$$

- 1. Find the optimal solution of this problem.
- 2. Write and solve the (Lagrangian) dual problem. Is there a duality gap ?

Exercise 6 (Two-way partitionning). Let $W \in S_n$ be a symmetric matrix, consider the follow- $\in [M_2]$ problem.

- $\begin{array}{ll} (P) & \underset{x \in \mathbb{R}^n}{\operatorname{Min}} & x^\top W x \\ & s.t. & x_i^2 = 1 & \forall i \in [n] \end{array}$
- 1. Consider a set of n element that you want to partition in 2 subsets, with a cost $c_{i,j}$ if i and j are in the same set, and a cost $-c_{i,j}$ if they are in a different set. Justify that it can be solved by solving (P).

- 2. Is (P) a convex problem ?
- 3. Show that, for any $\lambda \in \mathbb{R}^n$ such that $W + \text{diag}(\lambda) \succeq 0$, we have $\text{val}(P) \geq -\sum \lambda_i$. Deduce a lower bound on val(P).

Exercise 7 (Linear SVM : duality). Consider the following problem (see : https://www. youtube.com/watch?v=IOetFPgsMUc for background)

$$\min_{w \in \mathbb{R}^{d}, b \in \mathbb{R}} \frac{1}{2} \|w\|^{2}$$
s.t. $y_{i}(w^{\top}x_{i}+b) \geq 1 \qquad \forall i \in [n]$
 $\eta_{i} \geq 0 \qquad \forall i \in [n]$

- 1. In which case can we guarantee strong duality ?
- 2. Write the dual of this optimization problem and express the optimal primal solution (w^{\sharp}, b^{\sharp}) in terms of the optimal dual solution.

Exercise 8. We consider the following problem.

$$\min_{x_1, x_2} \qquad x_1^2 + x_2^2 \tag{1}$$

s.t.
$$(x_1 - 1)^2 + (x_2 - 1)^2 \le 1$$
 (2)
 $(x_1 - 1)^2 + (x_2 + 1)^2 \le 1$ (3)

- 1. Classify this problem (After 5th course)
- 2. Find the optimal solution and value of this problem.
- 3. Write and solve the KKT equation for this problem.
- 4. Derive and solve the Lagrangian dual of this problem.
- 5. Do we have strong duality ? If yes, could we have known it from the start ? If not, can you comment on why ?