

Contribution to Decomposition Methods in Stochastic Optimization

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PhD Defense
2014, June 25



UNIVERSITÉ
— PARIS-EST

Multistage Stochastic Optimization: an Example

How to manage a chain of dam producing electricity from the turbine water to optimize the gain?

$$\mathbb{E} \left[\sum_{i=1}^N \sum_{t=0}^{T-1} L_t^i(\underbrace{\mathbf{x}_t^i}_{\text{state}}, \underbrace{\mathbf{u}_t^i}_{\text{control}}, \underbrace{\mathbf{w}_{t+1}}_{\text{noise}}) \right]$$

Constraints:

- **dynamics:**
 $\mathbf{x}_{t+1} = f_t(\mathbf{x}_t, \mathbf{u}_t, \mathbf{w}_{t+1}),$
- **nonanticipativity:**
 $\mathbf{u}_t \preceq \mathcal{F}_t,$
- **spatial coupling:**
 $\mathbf{z}_t^{i+1} = g_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}^i).$

Figures/three_dams_VL.png

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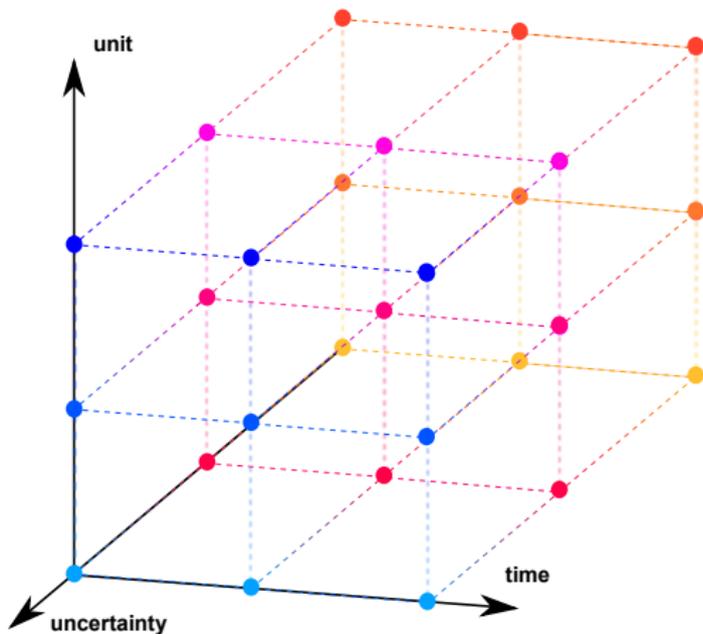
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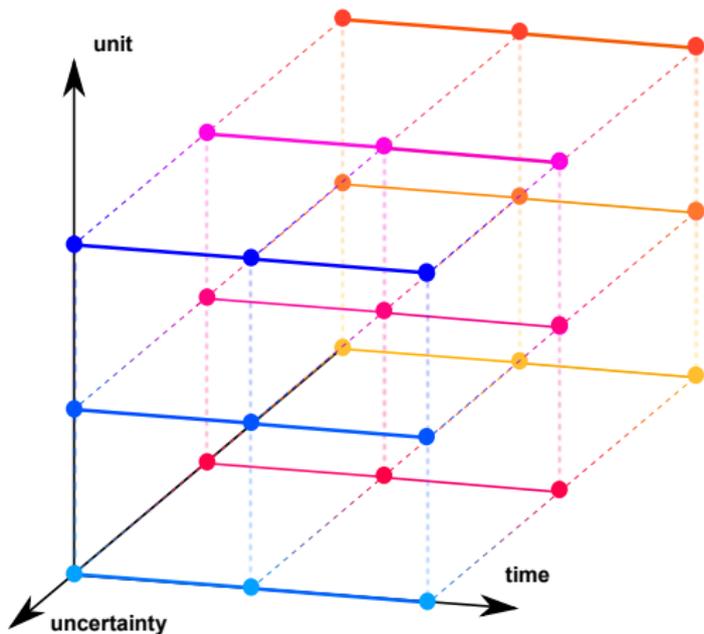
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Couplings for Stochastic Problems



$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

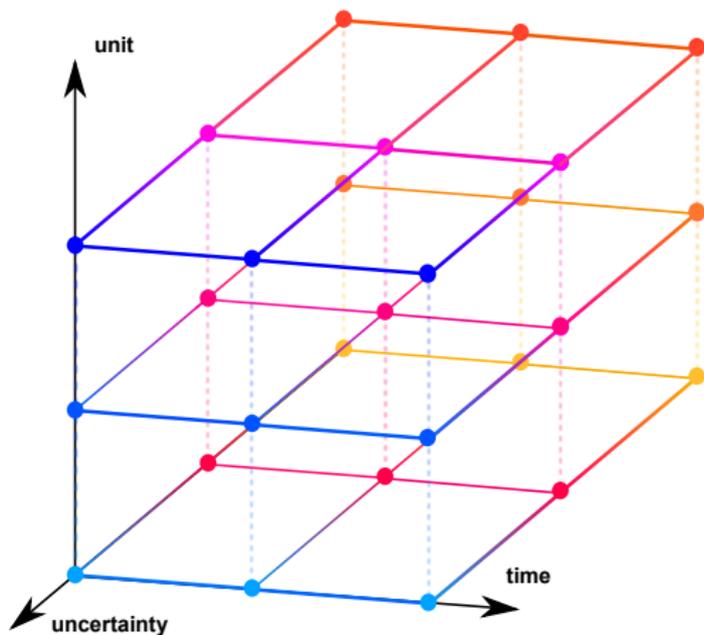
Couplings for Stochastic Problems: in Time



$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1})$$

$$\text{s.t. } \mathbf{X}_{t+1}^i = f_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1})$$

Couplings for Stochastic Problems: in Uncertainty

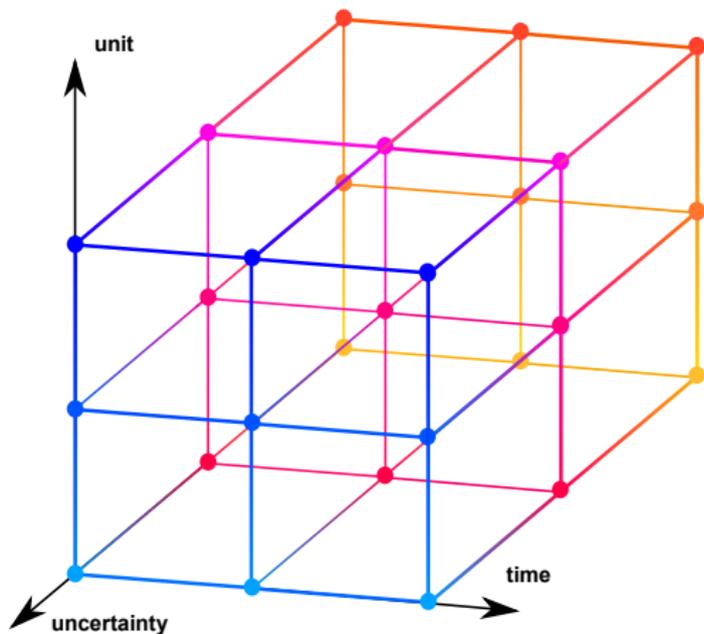


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$$\mathbf{U}_t^i \preceq \mathcal{F}_t = \sigma(\mathbf{W}_1, \dots, \mathbf{W}_t)$$

Couplings for Stochastic Problems: in Space



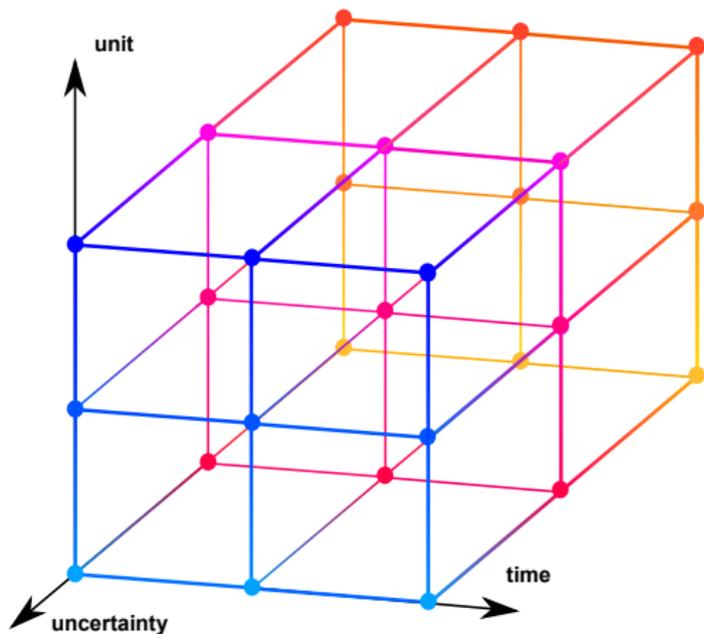
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$$\mathbf{u}_t^i \preceq \mathcal{F}_t = \sigma(\mathbf{w}_1, \dots, \mathbf{w}_t)$$

$$\sum_i \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) = 0$$

Couplings for Stochastic Problems: a Complex Problem



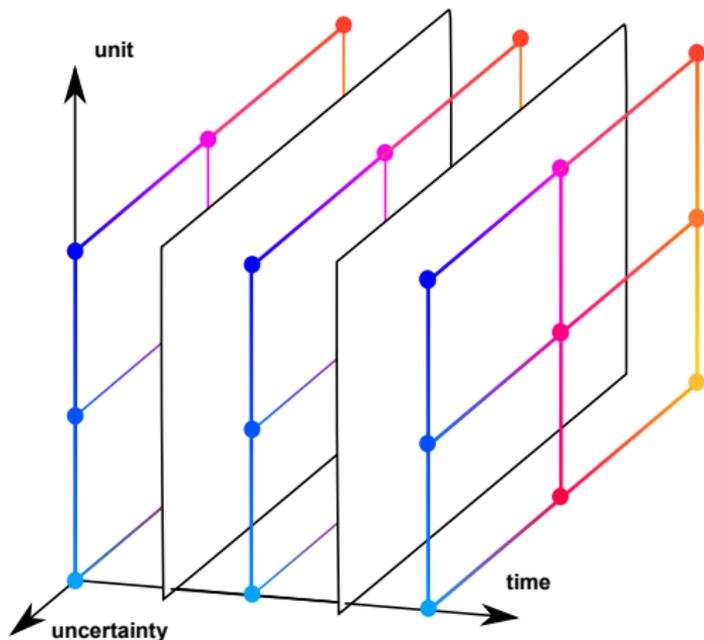
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Decompositions for Stochastic Problems: in Time



$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

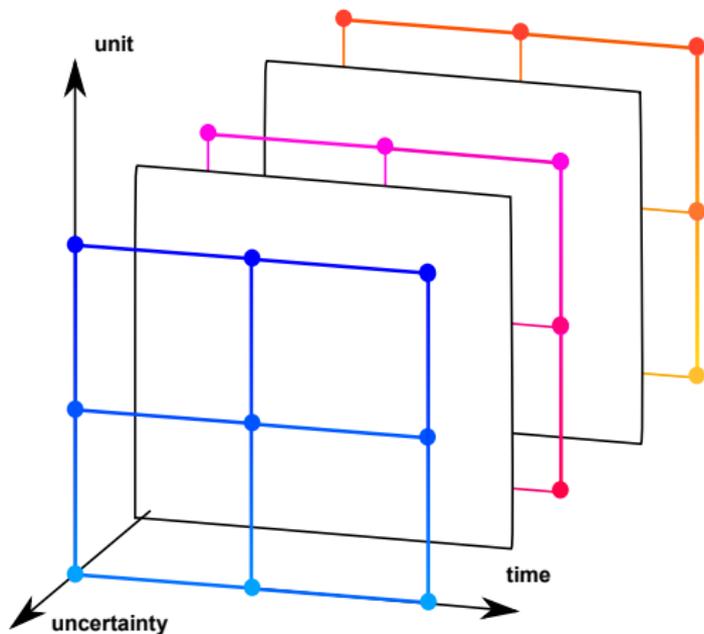
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Dynamic Programming
 Bellman (56)

Decompositions for Stochastic Problems: in Uncertainty



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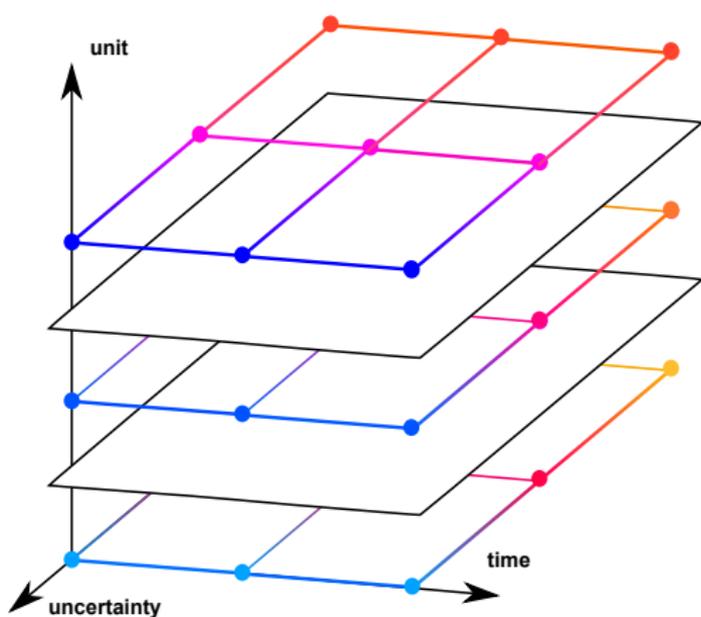
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Progressive Hedging
 Rockafellar - Wets (91)

Decompositions for Stochastic Problems: in Space



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Dual Approximate
 Dynamic Programming

Thesis Outline

- 1 Preliminaries
- 2 Time-Consistency: from Optimization to Risk Measures
- 3 Stochastic Dual Dynamic Programming Algorithm
- 4 Constraint Qualification in Stochastic Optimization
- 5 Constraint Qualification in (L^∞, L^1)
- 6 Uzawa Algorithm in L^∞
- 7 Epiconvergence of Relaxed Stochastic Problems
- 8 Dual Approximate Dynamic Programming Algorithm

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Presentation Outline

- 1 Time-Consistency: from Optimization to Risk Measures
 - A Framework for Dynamic Programming
 - Conditions for Time-Consistency
 - Examples
- 2 Spatial Stochastic Decomposition Method
 - Spatial Decomposition
 - Theoretical Results

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Classical Discrete Time Stochastic Optimization Problem

$$\begin{aligned} \min_{\mathbf{U}} \quad & \mathbb{E} \left[\overbrace{L_0(\mathbf{X}_0, \mathbf{U}_0, \mathbf{W}_1)}^{\text{instantaneous cost}} + \cdots + L_{T-1}(\mathbf{X}_{T-1}, \mathbf{U}_{T-1}, \mathbf{W}_T) + \overbrace{K(\mathbf{X}_T)}^{\text{final cost}} \right] \\ \text{s.t.} \quad & \mathbf{X}_0 = x_0 \\ & \mathbf{X}_{t+1} = f_t(\mathbf{X}_t, \mathbf{U}_t, \mathbf{W}_{t+1}) \quad (\text{dynamic}) \\ & \mathbf{U}_t \preceq \sigma(\mathbf{W}_1, \dots, \mathbf{W}_t) \quad (\text{non-anticipativity}) \end{aligned}$$

- \mathbf{X}_t : state (r.v. with value in \mathbb{X}_t),
- \mathbf{U}_t : control (r.v. with value in \mathbb{U}_t),
- \mathbf{W}_t : uncertainty (r.v. with value in \mathbb{W}_t)
 \rightsquigarrow time independence assumption!

Classical Discrete Time Stochastic Optimization Problem

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Risk Measure Formulation

$$\begin{aligned}
 \min_{\pi} \quad & \underbrace{q_{0,T}}_{\text{instantaneous cost}} \left\{ \underbrace{L_0(\mathbf{X}_0, \mathbf{U}_0, \mathbf{W}_1)}_{\text{instantaneous cost}}, \dots, L_{T-1}(\mathbf{X}_{T-1}, \mathbf{U}_{T-1}, \mathbf{W}_T), \underbrace{K(\mathbf{X}_T)}_{\text{final cost}} \right\} \\
 \text{s.t.} \quad & \mathbf{X}_0 = x_0 \\
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 & \mathbf{U}_t = \pi_t(\mathbf{X}_t) \quad \text{(non-anticipativity)}
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Example of Conditional Risk Measures

- $\varrho_{0,T}\{\mathbf{C}_0, \dots, \mathbf{C}_T\} = \mathbb{E}\left[\sum_{t=0}^T \mathbf{C}_t\right]$ (Classical framework)
- $\varrho_{0,T}\{\mathbf{C}_0, \dots, \mathbf{C}_T\} = \mathbb{E}\left[\sum_{t=0}^T r^t \mathbf{C}_t\right]$
- $\varrho_{0,T}\{\mathbf{C}_0, \dots, \mathbf{C}_T\} = \sup_{\mathbb{P} \in \mathcal{P}} \left\{ \mathbb{E}_{\mathbb{P}} \left[\sum_{t=0}^T \mathbf{C}_t \right] \right\}$
- $\varrho_{0,T}\{\mathbf{C}_0, \dots, \mathbf{C}_T\} = \sup_{\mathbb{P} \in \mathcal{P}} \left\{ \mathbb{E}_{\mathbb{P}} \left[\prod_{t=0}^T \mathbf{C}_t \right] \right\}$

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Dynamic Programming: Classical Framework

The sequence of Bellman functions $(V_t)_{t \in \llbracket 0, T \rrbracket}$ defined by

$$V_t(x) = \min_{\pi \in \Pi} \mathbb{E} \left[\sum_{\tau=t}^{T-1} L_\tau(\mathbf{X}_\tau, \pi_\tau(\mathbf{X}_\tau), \mathbf{W}_{\tau+1}) + K(\mathbf{X}_T) \right]$$

s.t. $\mathbf{X}_t = x$
 $\mathbf{X}_{\tau+1} = f_\tau(\mathbf{X}_\tau, \pi_\tau(\mathbf{X}_\tau), \mathbf{W}_{\tau+1})$

satisfies the Bellman equation \rightsquigarrow **Time Decomposition!**

$$\begin{cases} V_T(x) &= K(x) \\ V_t(x) &= \min_{u \in \mathcal{U}_t} \mathbb{E} \left[L_t(x, u, \mathbf{W}_{t+1}) + V_{t+1} \circ f_t(x, u, \mathbf{W}_{t+1}) \right] \end{cases}$$

Question: what about other risk measures?

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Uncertainty Aggregators

- Global uncertainty aggregator $\mathbb{G} : \mathcal{L}(\mathbb{W}_1 \times \dots \times \mathbb{W}_T; \mathbb{R}) \rightarrow \mathbb{R}$

- $\mathbb{G}[f] = \mathbb{E}[f(\mathbf{W}_1, \dots, \mathbf{W}_T)]$
- $\mathbb{G}[f] = \max_{w \in \mathbb{W}_1 \times \dots \times \mathbb{W}_T} f(w_1, \dots, w_T)$

- Time-step uncertainty aggregator $\mathbb{G}_t : \mathcal{L}(\mathbb{W}_t; \mathbb{R}) \rightarrow \mathbb{R}$

- $\mathbb{G}_t[f_t] = \mathbb{E}[f_t(\mathbf{W}_t)]$
- $\mathbb{G}_t[f_t] = \max_{w_t \in \mathbb{W}_t} f_t(w_t)$

- Composition of aggregators: $\mathbb{G}_t[w_t \mapsto \mathbb{G}_{t+1}[f(w_t, w_{t+1})]]$

$$\begin{aligned} & \max_{w \in \mathbb{W}_1 \times \dots \times \mathbb{W}_T} f(w_1, \dots, w_T) \\ &= \max_{w_1} \left[\max_{w_2} \left[\dots \max_{w_T} [f(w_1, \dots, w_T)] \right] \right] \end{aligned}$$

Time Aggregators

- Global time aggregator $\Phi : \mathbb{R}^{T+1} \rightarrow \mathbb{R}$
 - $\Phi\{c_0, \dots, c_T\} = \sum_{t=0}^T c_t$
 - $\Phi\{c_0, \dots, c_T\} = \prod_{t=0}^T c_t$
- Time-step time aggregator $\Phi_t : \mathbb{R}^2 \rightarrow \mathbb{R}$
 - $\Phi_t\{c_1, c_2\} = c_1 + c_2$
 - $\Phi_t\{c_1, c_2\} = c_1 \times c_2$
- Composition of aggregators $\Phi_t\{c_t, \Phi_{t+1}\{c_{t+1}, c_{t+2}\}\}$

$$\sum_{t=0}^T c_t = c_0 + \left\{ c_1 + \left\{ \dots + \{c_{T-1} + c_T\} \right\} \right\}$$

Constructing Optimization Problems

Time - then - Uncertainty (TU)

$$\varrho_{0,T}(\mathbf{C}_0, \dots, \mathbf{C}_T) = \mathbb{G} \left[\Phi \left\{ \mathbf{C}_0, \dots, \mathbf{C}_T \right\} \right]$$

Uncertainty - then - Time (UT)

$$\varrho_{0,T}(\mathbf{C}_0, \dots, \mathbf{C}_T) = \Phi \left\{ \mathbb{G}_0[\mathbf{C}_0], \dots, \mathbb{G}_T[\mathbf{C}_T] \right\}$$

(TU) examples:

- $\mathbb{E} \left[\sum_{t=0}^T \mathbf{C}_t \right]$
- $\mathbb{E} \left[\sum_{t=0}^T r^t \mathbf{C}_t \right]$
- $\max_{\mathbb{P} \in \mathcal{P}} \mathbb{E}_{\mathbb{P}} \left[\sum_{t=0}^T r^t \mathbf{C}_t \right]$

(UT) examples:

- $\sum_{t=0}^T \mathbb{E}[\mathbf{C}_t]$
- $\sum_{t=0}^T r^t \mathbb{E}[\mathbf{C}_t]$
- $\sum_{t=0}^T r^t \max_{\mathbb{P}_t \in \mathcal{P}_t} \mathbb{E}_{\mathbb{P}_t}[\mathbf{C}_t]$

Constructing Optimization Problems

Nested - Time - then - Uncertainty (NTU)

$$\varrho_{T,T}(\mathbf{C}_T) = \mathbb{G}_0 \left[\Phi_0 \left\{ \mathbf{C}_0, \mathbb{G}_1 \left[\Phi_1 \left\{ \dots \right. \right. \right. \right. \right. \\ \left. \left. \left. \left. \mathbb{G}_{T-1} \left[\Phi_{T-1} \left\{ \mathbf{C}_{T-1}, \mathbb{G}_T \left[\mathbf{C}_T \right] \right\} \right] \dots \right\} \right\} \right\} \right]$$

Nested - Uncertainty - then - Time (NUT)

$$\varrho_{0,T}(\mathbf{C}_0, \dots, \mathbf{C}_T) = \Phi_0 \left\{ \mathbb{G}_0 \left[\mathbf{C}_0 \right], \Phi_1 \left\{ \mathbb{G}_1 \left[\mathbf{C}_1 \right], \mathbb{G}_1 \left[\dots \right. \right. \right. \\ \left. \left. \left. \left. \Phi_{T-1} \left\{ \mathbb{G}_{T-1} \left[\mathbf{C}_{T-1} \right], \mathbb{G}_{T-1} \left[\mathbb{G}_T \left[\mathbf{C}_T \right] \right] \right\} \dots \right\} \right\} \right\}$$

Constructing Optimization Problems

Nested - Time - then - Uncertainty (NTU)

$$\varrho_{T-1,T}(\mathbf{C}_{T-1}, \mathbf{C}_T) = \mathbb{G}_0 \left[\Phi_0 \left\{ \mathbf{C}_0, \mathbb{G}_1 \left[\Phi_1 \left\{ \dots \right. \right. \right. \right. \right. \\ \left. \left. \left. \left. \mathbb{G}_{T-1} \left[\Phi_{T-1} \left\{ \mathbf{C}_{T-1}, \mathbb{G}_T \left[\mathbf{C}_T \right] \right\} \right] \dots \right\} \right\} \right\} \right]$$

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Constructing Optimization Problems

Nested - Time - then - Uncertainty (NTU)

$$\varrho_{1,T}(\mathbf{c}_1, \dots, \mathbf{c}_T) = \mathbb{G}_0 \left[\Phi_0 \left\{ \mathbf{c}_0, \mathbb{G}_1 \left[\Phi_1 \left\{ \dots \right. \right. \right. \right. \right. \\ \left. \left. \left. \left. \mathbb{G}_{T-1} \left[\Phi_{T-1} \left\{ \mathbf{c}_{T-1}, \mathbb{G}_T \left[\mathbf{c}_T \right] \right\} \right] \dots \right\} \right\} \right\} \right]$$

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Nested - Uncertainty - then - Time (NUT)

$$\varrho_{0,T}(\mathbf{C}_0, \dots, \mathbf{C}_T) = \Phi_0 \left\{ \mathbb{G}_0 [\mathbf{C}_0], \Phi_1 \left\{ \mathbb{G}_1 [\mathbf{C}_1], \mathbb{G}_1 \left[\dots \right. \right. \right. \\ \left. \left. \left. \Phi_{T-1} \left\{ \mathbb{G}_{T-1} [\mathbf{C}_{T-1}], \mathbb{G}_{T-1} [\mathbb{G}_T [\mathbf{C}_T]] \right\} \dots \right\} \right\} \right\}$$

Constructing Optimization Problems

Nested - Time - then - Uncertainty (NTU)

$$\varrho_{0,T}(\mathbf{c}_0, \dots, \mathbf{c}_T) = \mathbb{G}_0 \left[\Phi_0 \left\{ \mathbf{c}_0, \mathbb{G}_1 \left[\Phi_1 \left\{ \dots \right. \right. \right. \right. \right. \\ \left. \left. \left. \left. \mathbb{G}_{T-1} \left[\Phi_{T-1} \left\{ \mathbf{c}_{T-1}, \mathbb{G}_T [\mathbf{c}_T] \right\} \right] \dots \right\} \right\} \right\} \right]$$

Nested - Uncertainty - then - Time (NUT)

$$\varrho_{0,T}(\mathbf{c}_0, \dots, \mathbf{c}_T) = \Phi_0 \left\{ \mathbb{G}_0 [\mathbf{c}_0], \Phi_1 \left\{ \mathbb{G}_1 [\mathbf{c}_1], \mathbb{G}_1 \left[\dots \right. \right. \right. \right. \\ \left. \left. \left. \left. \Phi_{T-1} \left\{ \mathbb{G}_{T-1} [\mathbf{c}_{T-1}], \mathbb{G}_{T-1} [\mathbb{G}_T [\mathbf{c}_T]] \right\} \dots \right\} \right\} \right\}$$

Conditions for a Dynamic Programming Principle (NTU)

De Lara - L.

Assume that the time-step aggregators \mathbb{G}_t and Φ_t are **monotonous**. Define the value functions

$$\begin{cases} V_T^{\text{NTU}}(x) &= K(x) \\ V_t^{\text{NTU}}(x) &= \inf_{u \in \mathbb{U}_t} \mathbb{G}_t \left[\Phi_t \left\{ L_t(x, u, \cdot), V_{t+1}^{\text{NTU}} \circ f_t(x, u, \cdot) \right\} \right] \end{cases}$$

Assume that there exists an admissible strategy π^\sharp such that

$$\pi_t^\sharp(x) \in \arg \min_{u \in \mathbb{U}_t} \mathbb{G}_t \left[\Phi_t \left\{ L_t(x, u, \cdot), V_{t+1}^{\text{NTU}} \circ f_t(x, u, \cdot) \right\} \right]$$

Then, π^\sharp is an **optimal policy**.

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$$\begin{cases} V_7^{\text{NUT}}(x) &= K(x) \\ V_t^{\text{NUT}}(x) &= \inf_{u \in \mathbb{U}_t} \Phi_t \left\{ \mathbb{G}_t [L_t(x, u, \cdot)], \mathbb{G}_t [V_{t+1}^{\text{NUT}} \circ f_t(x, u, \cdot)] \right\} \end{cases}$$

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Then, π^\sharp is an **optimal policy**.

Commutation

Commutation

Uncertainty aggregator \mathbb{G}_{t+1} and time-aggregator Φ_t are said to be commuting when, for all functions f and g

$$\mathbb{G}_{t+1} \left[\Phi_t \{ f(\mathbf{W}_t), g(\mathbf{W}_{t+1}) \} \right] = \Phi_t \left\{ f(\mathbf{W}_t), \mathbb{G}_{t+1} [g(\mathbf{W}_{t+1})] \right\}$$

Examples:

- $\mathbb{E}_{\mathbb{P}_{t+1}} [f(\mathbf{W}_t) + g(\mathbf{W}_{t+1})] = f(\mathbf{W}_t) + \mathbb{E}_{\mathbb{P}_{t+1}} [g(\mathbf{W}_{t+1})]$
- commutation with sum \iff translation equivariance property
- $\mathbb{E}_{\mathbb{P}_{t+1}} [f(\mathbf{W}_t) \times g(\mathbf{W}_{t+1})] = f(\mathbf{W}_t) \times \mathbb{E}_{\mathbb{P}_{t+1}} [g(\mathbf{W}_{t+1})]$

Conditions for a Dynamic Programming Principle (TU)

De Lara - L.

Assume that

- the **global aggregators are a composition** of time-step aggregators,
- the time-step aggregators \mathbb{G}_t and Φ_t are **monotonous**,
- the time-step aggregators \mathbb{G}_t and Φ_s ($s < t$) **commute**.

Then, the **nested and not nested formulations are equivalent**, and **we have a DP equation**.

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Time-Consistency of a Sequence of Optimization Problems

$$\begin{aligned} (\mathcal{P}_t) \quad & \min_{\pi} \varrho_{t,T} \left(L_t(\mathbf{X}_t, \mathbf{U}_t, \mathbf{W}_{t+1}), \dots, \right. \\ & \left. L_{T-1}(\mathbf{X}_{T-1}, \mathbf{U}_{T-1}, \mathbf{W}_T), K(\mathbf{X}_T) \right) \\ \text{s.t.} \quad & \mathbf{X}_t = \mathbf{x} \\ & \mathbf{X}_{\tau+1} = f_{\tau}(\mathbf{X}_{\tau}, \mathbf{U}_{\tau}, \mathbf{W}_{\tau}) \\ & \mathbf{U}_{\tau} = \pi_{\tau}(\mathbf{X}_{\tau}) \end{aligned}$$

The sequence of problems $(\mathcal{P}_t)_{t \in \llbracket 0, T-1 \rrbracket}$ is said to be **time consistent** if there exists an optimal strategy of Problem (\mathcal{P}_{t_0}) such that its restriction is optimal for (\mathcal{P}_{t_1}) , $(t_1 > t_0)$.

Time-Consistency of a Dynamic Risk Measure

A sequence of conditional risk measures $(\rho_{0,T}, \rho_{1,T}, \dots, \rho_T)$ is **time-consistent** if for any two sequences of costs $(\mathbf{C}_0, \dots, \mathbf{C}_T)$ $(\mathbf{C}'_0, \dots, \mathbf{C}'_T)$ we have

$$\left. \begin{aligned} (\mathbf{C}_{t_1}, \dots, \mathbf{C}_{t_2-1}) &= (\mathbf{C}'_{t_1}, \dots, \mathbf{C}'_{t_2-1}) \\ \rho_{t_2,T}(\mathbf{C}_{t_2}, \dots, \mathbf{C}_T) &\leq \rho_{t_2,T}(\mathbf{C}'_{t_2}, \dots, \mathbf{C}'_T) \end{aligned} \right\}$$

$$\implies \rho_{t_1,T}(\mathbf{C}_{t_1}, \dots, \mathbf{C}_{t_2}, \dots, \mathbf{C}_T) \leq \rho_{t_1,T}(\mathbf{C}'_{t_1}, \dots, \mathbf{C}'_{t_2}, \dots, \mathbf{C}'_T)$$

Time-Consistency Result

Nested formulation - De Lara, L.

If the time-step aggregators are **monotonous**, the induced:

- sequence of optimization problems
- sequence of conditional risk measures

are **time consistent**.

Non-Nested Formulation - De Lara, L.

If the global aggregators are **composition** of **monotonous** and **commuting** time-step aggregators, the induced

- sequence of optimization problems
- sequence of conditional risk measures

are **time consistent**.

Markovian Case

- We have extended the framework to allow for Markovian aggregators:

$$\mathbb{G}_t \rightsquigarrow \mathbb{G}_t^x \quad \Phi_t \rightsquigarrow \Phi_t^x$$

- Examples:
 - Conditional expectation: $\mathbb{G}_t^x = \mathbb{E}[\cdot \mid \mathbf{X}_t = x]$,
 - Markov risk measure (Ruszczynski 2010).

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Classical Extension: Multiplicative Case

A stochastic viability problem can be written

$$\begin{aligned} \max_{\pi \in \Pi} \quad & \mathbb{P}\left(\left\{\mathbf{X}_t \in \mathcal{X}_t, \quad \forall t \in \llbracket 0, T \rrbracket\right\}\right) \\ \text{s.t.} \quad & \mathbf{X}_{t+1} = f_t(\mathbf{X}_t, \mathbf{U}_t, \mathbf{W}_{t+1}) \\ & \mathbf{U}_t = \pi_t(\mathbf{X}_t) \end{aligned}$$

With the following DP equation

$$\begin{cases} V_T(x) &= \mathbb{E}\left[\mathbb{1}_{\{x \in \mathcal{X}_T\}}\right] \\ V_t(x) &= \max_{u \in \mathbf{U}_t} \mathbb{E}\left[\mathbb{1}_{\{x \in \mathcal{X}_t\}} \cdot V_{t+1} \circ f_t(x, u, \mathbf{W}_{t+1})\right] \end{cases}$$

Classical Extension: Multiplicative Case

A stochastic viability problem can be written

$$\begin{aligned} \max_{\pi \in \Pi} \quad & \mathbb{E} \left[\prod_{t=0}^T \mathbb{1}_{\{\mathbf{X}_t \in \mathcal{X}_t\}} \right] \\ \text{s.t.} \quad & \mathbf{X}_{t+1} = f_t(\mathbf{X}_t, \mathbf{U}_t, \mathbf{W}_{t+1}) \\ & \mathbf{U}_t = \pi_t(\mathbf{X}_t) \end{aligned}$$

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Coherent Risk Measure

Consider the following sequence of conditional risk measures.

$$\varrho_{t,T}(\mathbf{C}) = \sup_{\mathbb{P}_t \in \mathcal{P}_t} \mathbb{E}_{\mathbb{P}_t} \left[\cdots \sup_{\mathbb{P}_T \in \mathcal{P}_T} \mathbb{E}_{\mathbb{P}_T} \left[\sum_{s=t}^T \left(\alpha_s(\mathbf{c}_s) \prod_{r=t}^{s-1} \beta_r(\mathbf{c}_r) \right) \right] \cdots \right]$$

The associated optimization problem is solved by the following DP equation (if $\beta_t \geq 0$)

$$\begin{cases} V_T(x) &= K(x) \\ V_t(x) &= \inf_u \sup_{\mathbb{P}_t \in \mathcal{P}_t} \left\{ \mathbb{E}_{\mathbb{P}_t} \left[\alpha_t(L_t(x, u, \cdot)) + \beta_t(L_t(x, u, \cdot)) V_{t+1} \circ f_t(x, u, \cdot) \right] \right\} \end{cases}$$

Elements of proof

- The problem is of (TU) form where the global aggregators are **composition** of the following time-step aggregators:

$$\begin{cases} \mathbb{G}_t[\cdot] &= \sup_{\mathbb{P}_t \in \mathcal{P}_t} \mathbb{E}_{\mathbb{P}_t}[\cdot] \\ \Phi_t\{c, c'\} &= \alpha_t(c) + \beta_t(c)c' \end{cases}$$

- The time-step aggregators are **monotonous**.
- The time-step aggregators **commute**:

$$\begin{aligned} \mathbb{G}_t[\Phi_s\{\mathbf{C}_s, \mathbf{C}_t\}] &= \sup_{\mathbb{P}_t \in \mathcal{P}_t} \left(\mathbb{E}_{\mathbb{P}_t}[\alpha_s(\mathbf{C}_s) + \beta_s(\mathbf{C}_s)\mathbf{C}_t] \right) \\ &= \sup_{\mathbb{P}_t \in \mathcal{P}_t} \left(\alpha_s(\mathbf{C}_s) + \beta_s(\mathbf{C}_s)\mathbb{E}_{\mathbb{P}_t}[\mathbf{C}_t] \right) && \text{Translation-equiv.} \\ &= \alpha_s(\mathbf{C}_s) + \beta_s(\mathbf{C}_s) \sup_{\mathbb{P}_t \in \mathcal{P}_t} \left(\mathbb{E}_{\mathbb{P}_t}[\mathbf{C}_t] \right) && \text{Pos. Homogeneity} \\ &= \Phi_s\{\mathbf{C}_s, \mathbb{G}_t[\mathbf{C}_t]\} \end{aligned}$$

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Conclusion of Part I

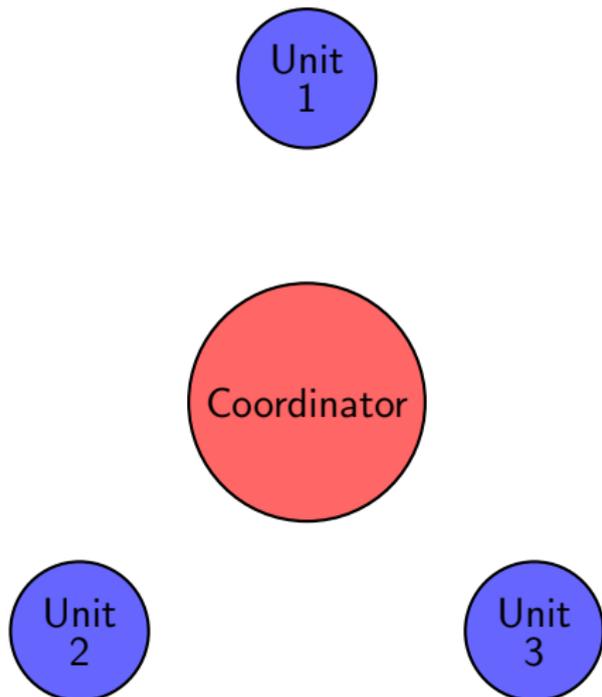
- We have presented a generic framework for stochastic optimization problem and conditions to write a chained time decomposition through a DP equation. We extended it to a Markovian framework.
- We show that our conditions lead to time-consistency of
 - the sequence of induced optimization problems,
 - and the induced dynamic risk measure.
- This part was concerned with formulation of problem in a time-consistent way, and time decomposition. However, it is still affected by the so-called “curse of dimensionality”.

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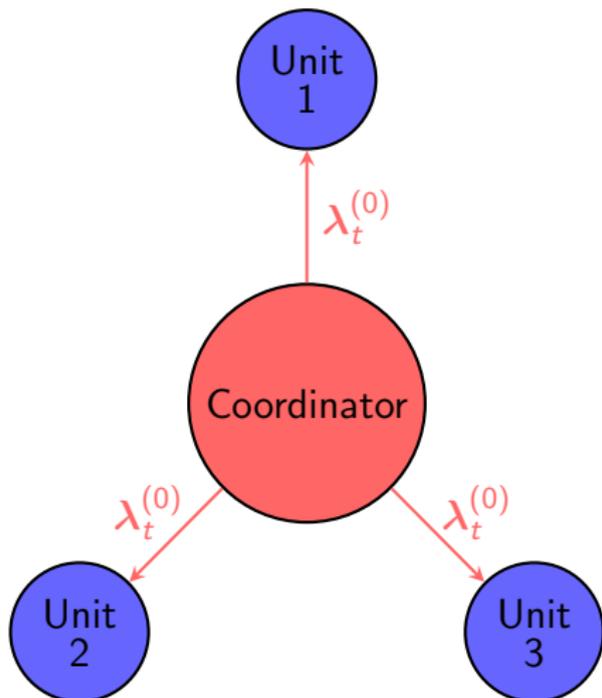
Intuition of Spatial Decomposition

- Satisfy a demand (over T time step) with N units of production at minimal cost.
- **Price decomposition:**
 - the coordinator sets a sequence of price λ_t ,
 - the units send their production planning $\mathbf{u}_t^{(i)}$,
 - the coordinator compares total production and demand and updates the price,
 - and so on...



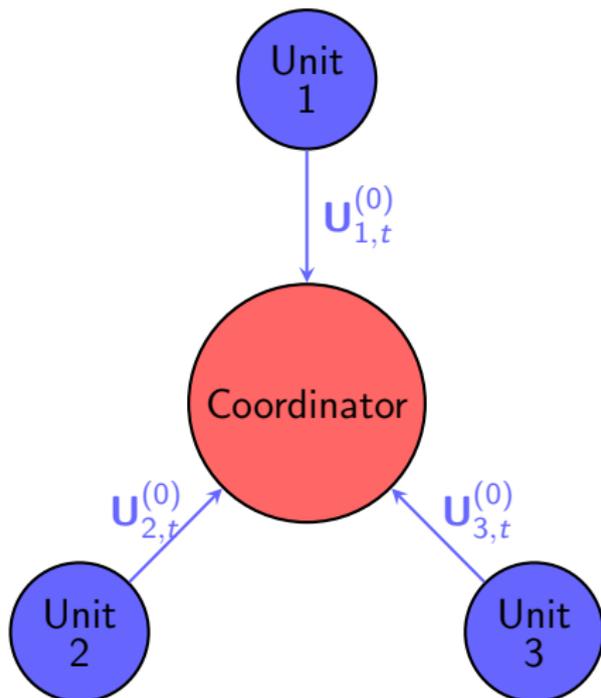
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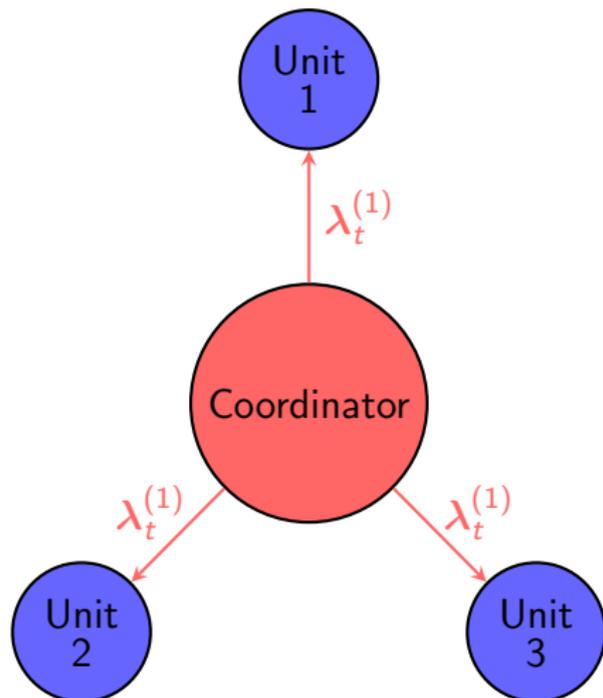
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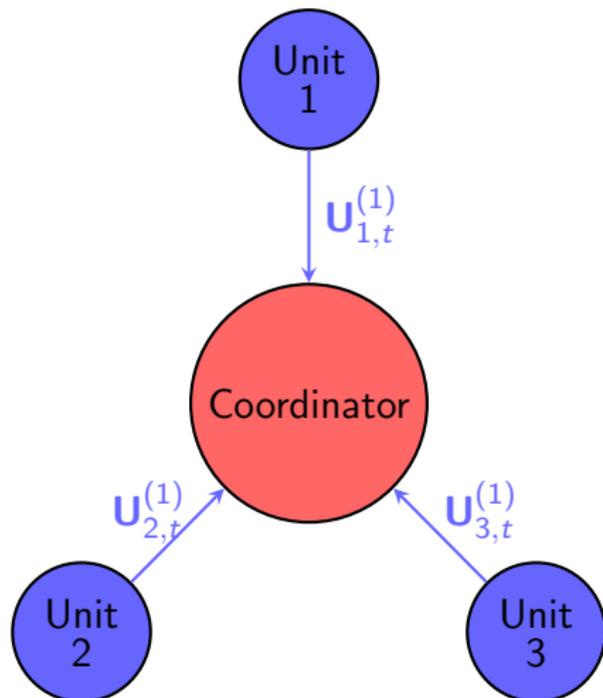
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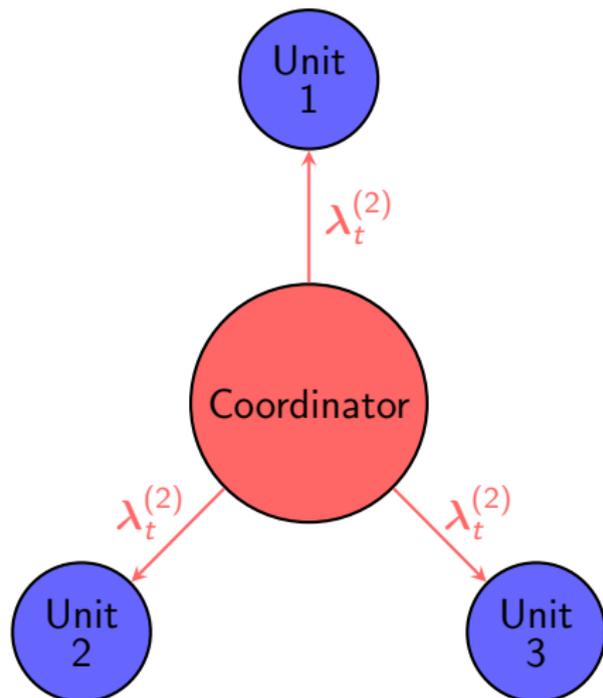
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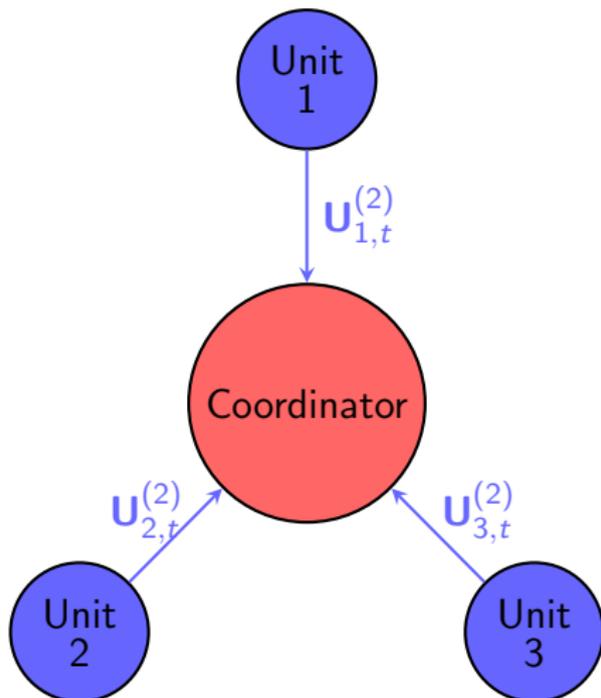
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Intuition of Spatial Decomposition

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 - and so on...



Primal Problem

$$\begin{aligned} \min_{\mathbf{X}, \mathbf{U}} \quad & \sum_{i=1}^N \mathbb{E} \left[\sum_{t=0}^T L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + K^i(\mathbf{x}_T^i) \right] \\ \forall i, \quad & \mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}), \quad \mathbf{x}_0^i = x_0^i, \\ \forall i, \quad & \mathbf{u}_t^i \in \mathcal{U}_{t,i}^{ad}, \quad \mathbf{u}_t^i \preceq \mathcal{F}_t, \\ & \sum_{i=1}^N \theta_t^i(\mathbf{u}_t^i) = 0 \end{aligned}$$

Solvable by DP with state $(\mathbf{X}_1, \dots, \mathbf{X}_N)$ (under noise independence assumption)

Primal Problem

$$\begin{aligned}
 \min_{\mathbf{X}, \mathbf{U}} \quad & \sum_{i=1}^N \mathbb{E} \left[\sum_{t=0}^T L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + K^i(\mathbf{x}_T^i) \right] \\
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 \sum_{i=1}^N \theta_t^i(\mathbf{u}_t^i) = 0 \quad & \rightsquigarrow \lambda_t \quad \text{multiplier}
 \end{aligned}$$

Solvable by DP with state $(\mathbf{X}_1, \dots, \mathbf{X}_N)$ (under noise independence assumption)

Primal Problem with Dualized Constraint

$$\begin{aligned} \min_{\mathbf{x}, \mathbf{U}} \max_{\lambda} \sum_{i=1}^N \mathbb{E} \left[\sum_{t=0}^T L_t^i(\mathbf{x}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}) + \langle \lambda_t, \theta_t^i(\mathbf{U}_t^i) \rangle + K^i(\mathbf{x}_T^i) \right] \\ \forall i, \quad \mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}), \quad \mathbf{x}_0^i = x_0^i, \\ \forall i, \quad \mathbf{U}_t^i \in \mathcal{U}_{t,i}^{ad}, \quad \mathbf{U}_t^i \preceq \mathcal{F}_t, \end{aligned}$$

Coupling constraint dualized \implies remaining constraints are i by i

Dual Problem

$$\begin{aligned} \max_{\lambda} \min_{\mathbf{X}, \mathbf{U}} \sum_{i=1}^N \mathbb{E} \left[\sum_{t=0}^T L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + \langle \lambda_t, \theta_t^i(\mathbf{u}_t^i) \rangle + K^i(\mathbf{x}_T^i) \right] \\ \forall i, \quad \mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}), \quad \mathbf{x}_0^i = x_0^i, \\ \forall i, \quad \mathbf{u}_t^i \in \mathcal{U}_{t,i}^{ad}, \quad \mathbf{u}_t^i \preceq \mathcal{F}_t, \end{aligned}$$

Exchange operator **min** and **max** to obtain a new problem

Decomposed Dual Problem

$$\max_{\lambda} \sum_{i=1}^N \min_{\mathbf{x}^i, \mathbf{u}^i} \mathbb{E} \left[\sum_{t=0}^T L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + \langle \lambda_t, \theta_t^i(\mathbf{u}_t^i) \rangle + K^i(\mathbf{x}_T^i) \right]$$
$$\mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}), \quad \mathbf{x}_0^i = x_0^i,$$
$$\mathbf{u}_t^i \in \mathcal{U}_{t,i}^{ad}, \quad \mathbf{u}_t^i \preceq \mathcal{F}_t,$$

For a given λ , minimum of sum is sum of minima

Inner Minimization Problem

$$\begin{aligned} \min_{\mathbf{x}^i, \mathbf{u}^i} \mathbb{E} & \left[\sum_{t=0}^T L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + \langle \lambda_t, \theta_t^i(\mathbf{u}_t^i) \rangle + K^i(\mathbf{x}_T^i) \right] \\ \mathbf{x}_{t+1}^i &= f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}), \quad \mathbf{x}_0^i = x_0^i, \\ \mathbf{u}_t^i &\in \mathcal{U}_{t,i}^{ad}, \quad \mathbf{u}_t^i \preceq \mathcal{F}_t, \end{aligned}$$

We have N smaller subproblems. Can they be solved by DP?

Inner Minimization Problem

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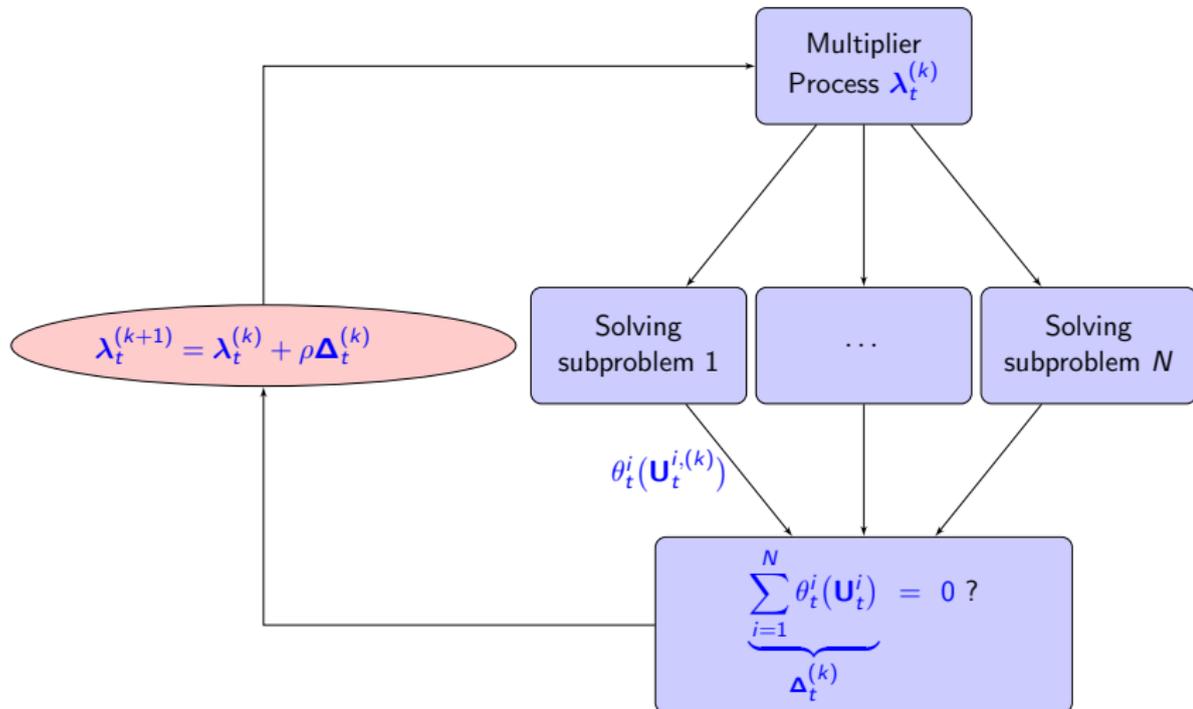
No: $\boldsymbol{\lambda}$ is a time-dependent noise $\rightsquigarrow \mathbf{x}_t^i$ is not a proper state, but rather $(\mathbf{w}_1, \dots, \mathbf{w}_t)$

A Few Questions

- What is the duality scheme ? In which space lives the multiplier process λ ?
 - L^2
 - L^1
 - $(L^\infty)^*$
- What are the relations between the primal and dual problem?
- Can we solve the subproblems by Dynamic Programming?
 \rightsquigarrow **No!** (with small enough state)
- How to update the multiplier process?
 \rightsquigarrow “gradient step”:

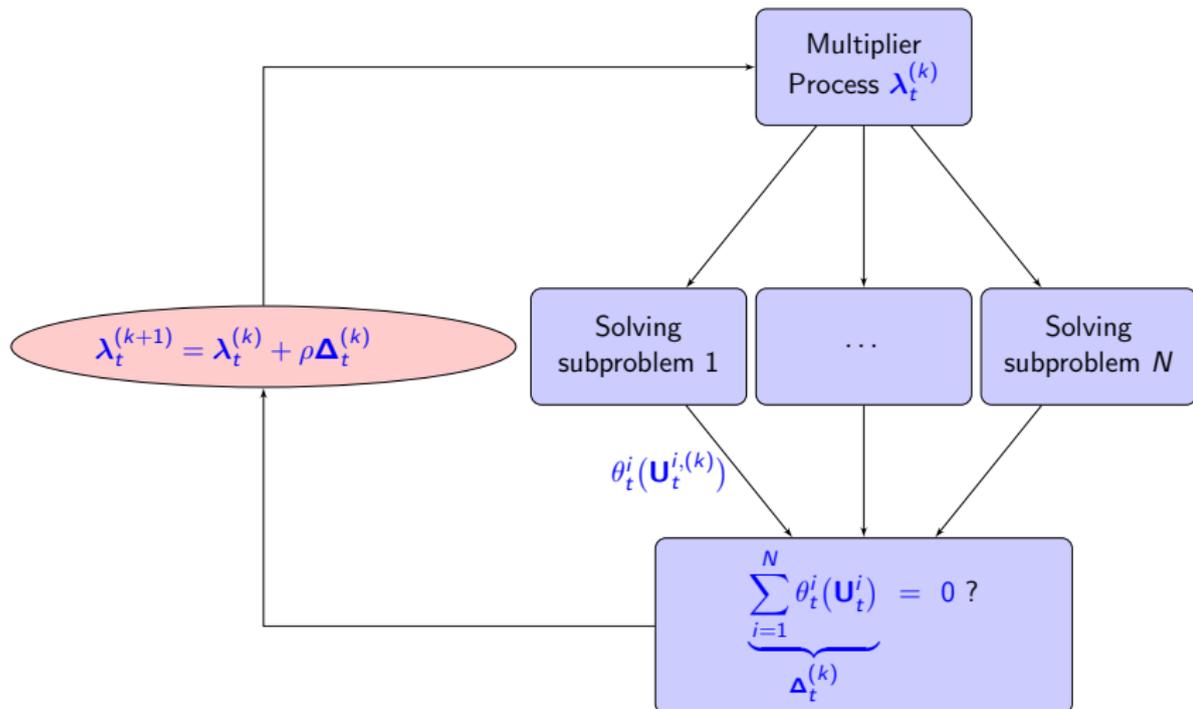
$$\lambda_t^{(k+1)} = \lambda_t^{(k)} + \rho \sum_{i=1}^N \theta_t^i(\mathbf{u}_t^{i,k})$$

Stochastic spatial decomposition scheme



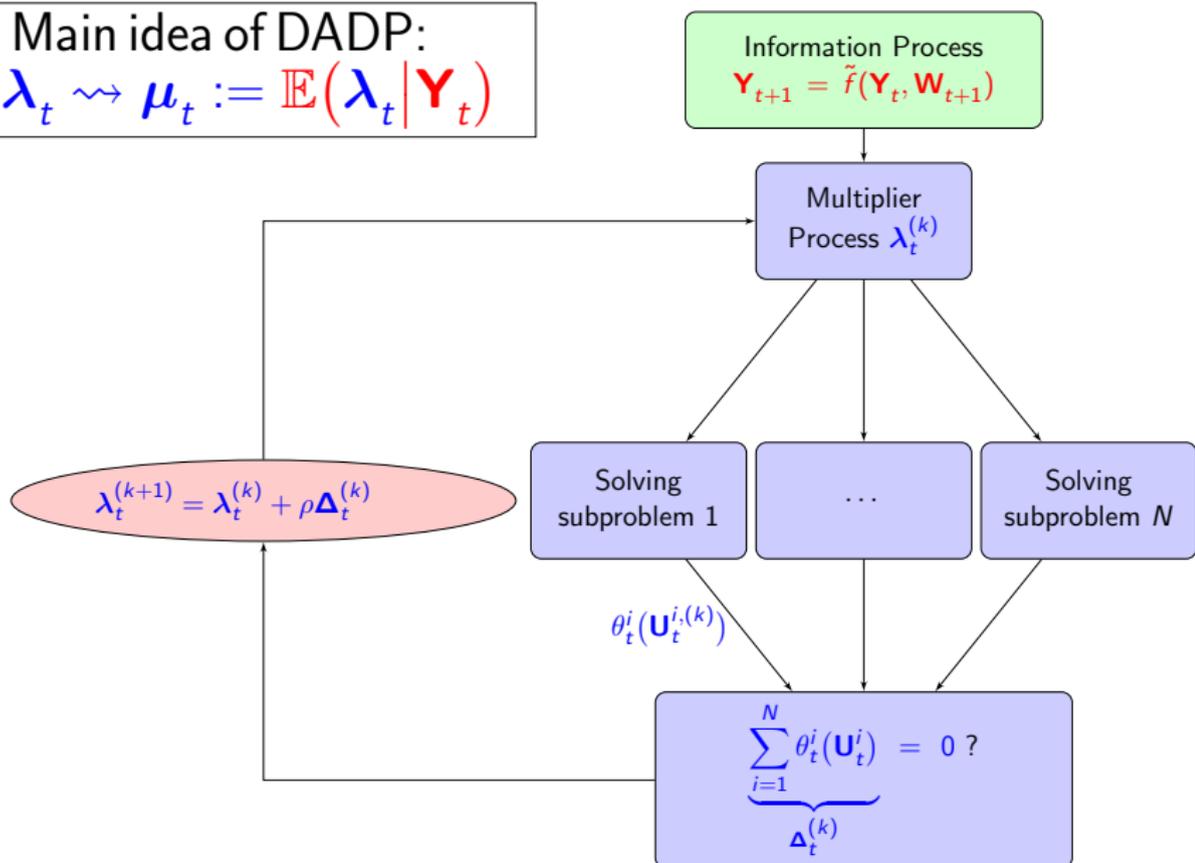
Main idea of DADP:

$$\lambda_t \rightsquigarrow \mu_t := \mathbb{E}(\lambda_t | \mathbf{Y}_t)$$



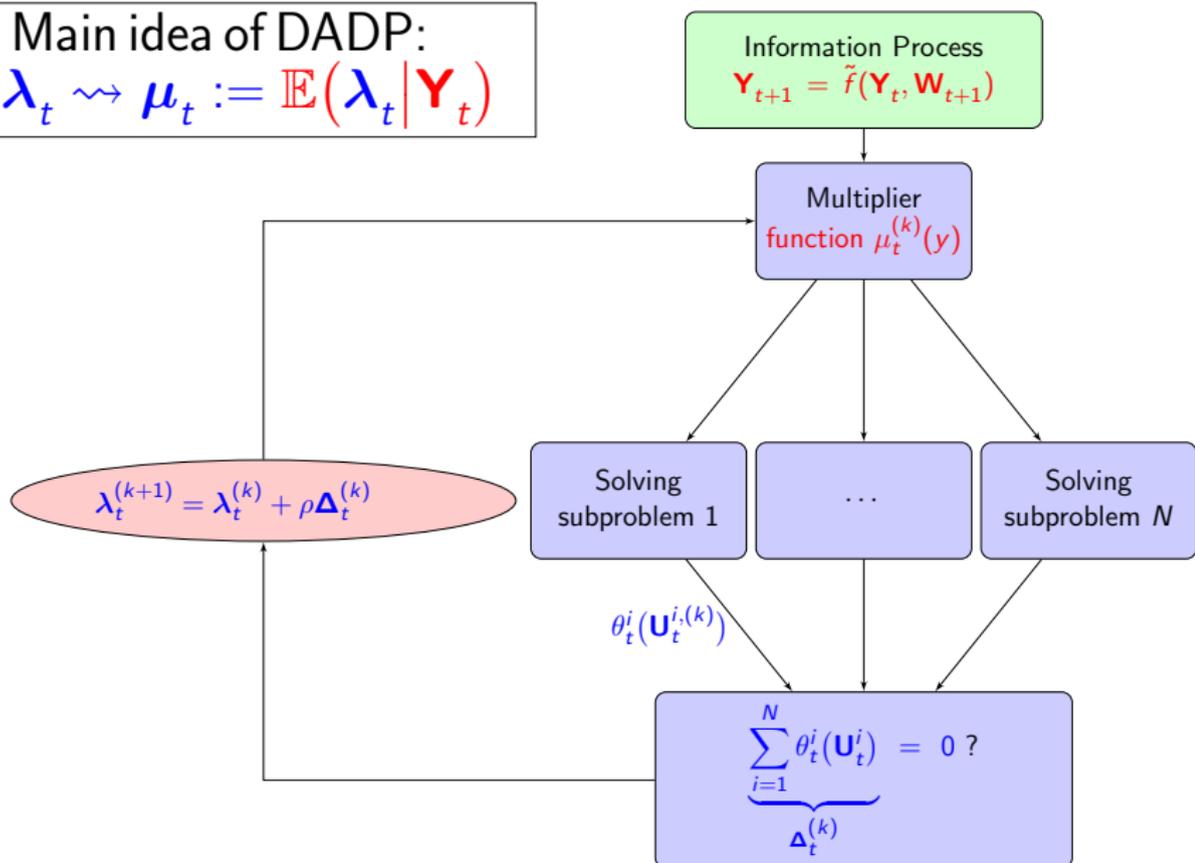
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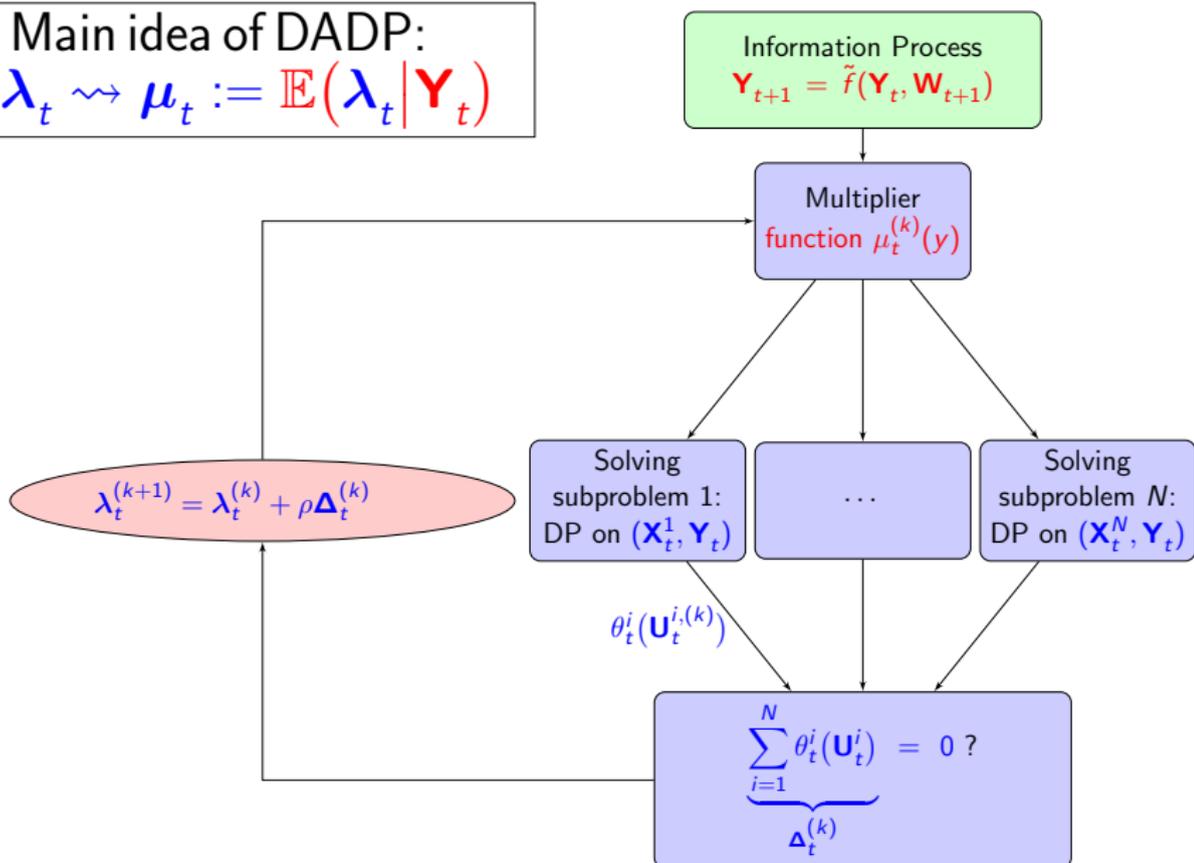
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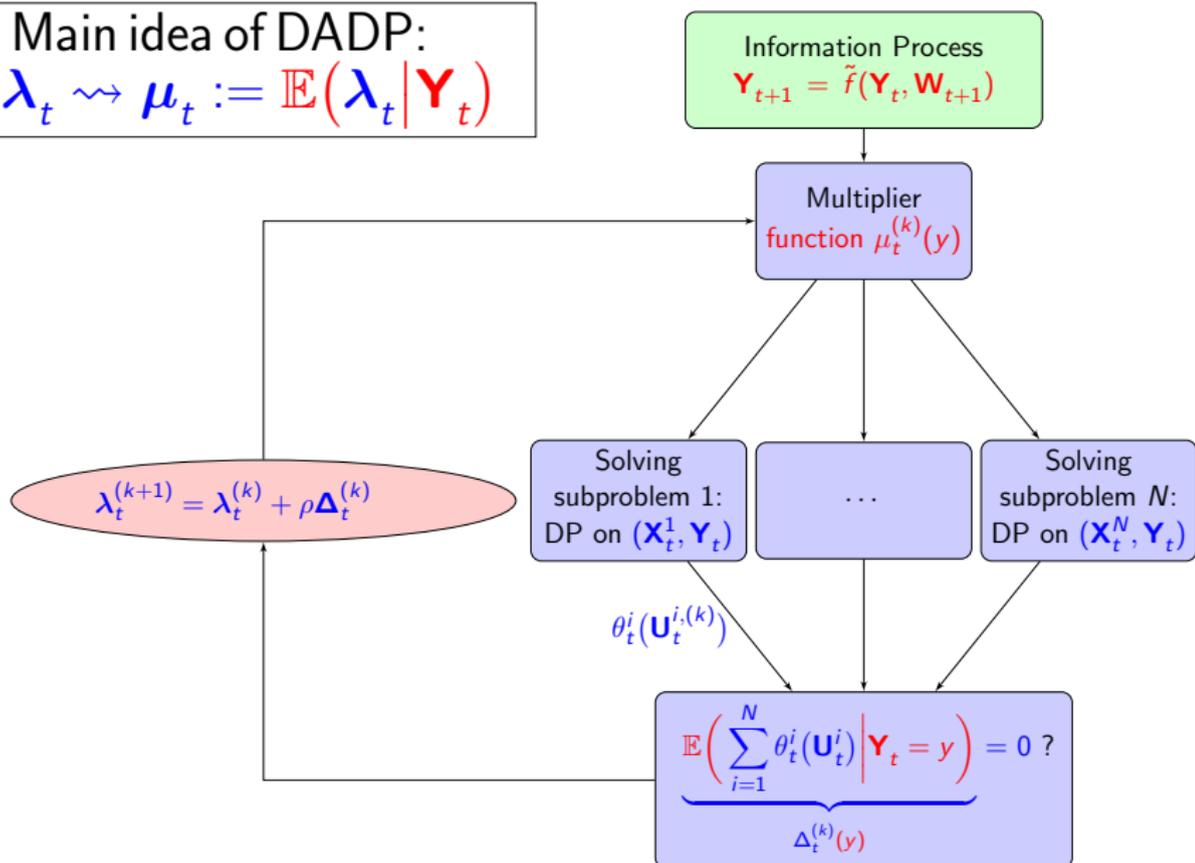
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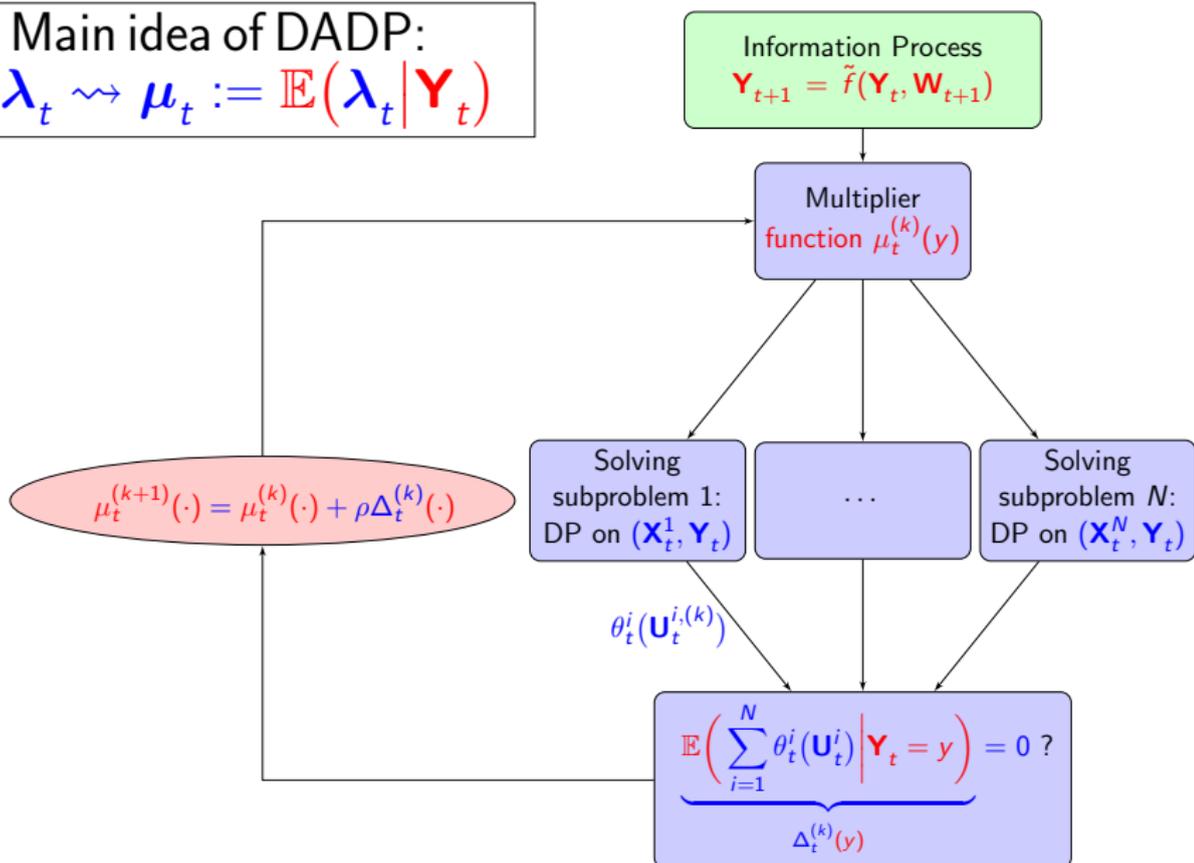
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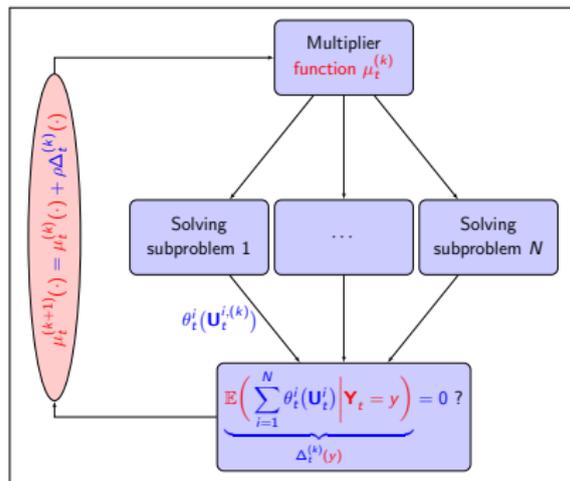
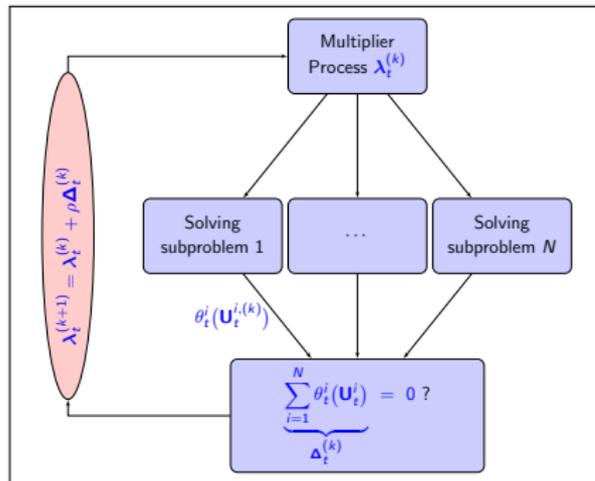


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Main idea of DADP: $\lambda_t \rightsquigarrow \mu_t := \mathbb{E}(\lambda_t | \mathbf{Y}_t)$



Main problems:

- Subproblems not easily solvable by DP
- $\lambda^{(k)}$ live in a huge space

Advantages:

- Subproblems solvable by DP with state $(\mathbf{X}_t^i, \mathbf{Y}_t)$
- $\mu^{(k)}$ live in a smaller space

Presentation Outline

- 1 Time-Consistency: from Optimization to Risk Measures
 - A Framework for Dynamic Programming
 - Conditions for Time-Consistency
 - Examples
- 2 Spatial Stochastic Decomposition Method
 - Spatial Decomposition
 - Theoretical Results

Three Interpretations of DADP

- DADP as an approximation of the optimal multiplier

$$\lambda_t \rightsquigarrow \mathbb{E}(\lambda_t | \mathbf{Y}_t) .$$

- DADP as a decision-rule approach in the dual

$$\max_{\lambda} \min_{\mathbf{U}} L(\lambda, \mathbf{U}) \rightsquigarrow \max_{\lambda_t \preceq \mathbf{Y}_t} \min_{\mathbf{U}} L(\lambda, \mathbf{U}) .$$

- DADP as a constraint relaxation in the primal

$$\sum_{i=1}^n \theta_t^i(\mathbf{u}_t^i) = 0 \rightsquigarrow \mathbb{E} \left(\sum_{i=1}^n \theta_t^i(\mathbf{u}_t^i) \middle| \mathbf{Y}_t \right) = 0 .$$

Consistence of the Approximation Scheme

- The DADP algorithm solves a relaxation (\mathcal{P}_Y) of the original problem (\mathcal{P}) where

$$\sum_{i=1}^n \theta_t^i(\mathbf{u}_t^i) = 0 \quad \rightsquigarrow \quad \mathbb{E} \left(\sum_{i=1}^n \theta_t^i(\mathbf{u}_t^i) \middle| \mathbf{Y}_t \right) = 0$$

- Question: if we consider a sequence of information processes $\{\mathbf{Y}^{(n)}\}_{n \in \mathbb{N}}$, such that the information converges

$$\sigma(\mathbf{Y}_t^{(n)}) \rightarrow \sigma(\mathbf{W}_0, \dots, \mathbf{W}_t)$$

does the associated sequence $(\mathbf{u}^{\mathbf{Y}^{(n)}})$ of optimal control converges toward an optimal control of (\mathcal{P}) ?

Epiconvergence of Approximation

Epiconvergence result -L.

Assume that

- the cost functions L_t^i , dynamic functions f_t^i and constraint functions θ_t^i are continuous;
- the noise variables \mathbf{W}_t are essentially bounded;
- the constraint sets $\mathcal{U}_{i,t}^{\text{ad}}$ are bounded.

Consider a sequence of information process $\{\mathbf{Y}^{(n)}\}_{n \in \mathbb{N}}$ such that $\sigma(\mathbf{Y}^{(n)}) \rightarrow \mathcal{F}_\infty$. Let $\mathbf{U}^{(n)}$ be an ε_n -optimal solution to the relaxed problem $(\mathcal{P}^{\mathbf{Y}^{(n)}})$.

Then, every cluster point^a of $\{\mathbf{U}^{(n)}\}_{n \in \mathbb{N}}$ is an optimal solution of the relaxation corresponding to \mathcal{F}_∞ .

^afor the topology of the convergence in probability

Convergence of Coordination Method

- We consider a given information process \mathbf{Y} .
- Question: does the algorithm
 - 1 solve the N subproblems
 - 2 update the multiplier by a gradient-stepyield a converging sequence of controls $\mathbf{U}^{(k)}$?
- It is an application of the so-called Uzawa algorithm. This algorithm take naturally place in an Hilbert space, here L^2 is the natural choice. However, existence of saddle-point in L^2 is difficult to prove. Hence we adapt the algorithm to a non-reflexive Banach space: L^∞ .

Coordination-Convergence Result

Convergence result - Carpentier, L.

Assume that,

- the set of uncertainties is finite;
- the local cost L_t^i are Gâteaux-differentiable functions, strongly convex (in (x, u)) and continuous (in w);
- the evolution functions f_t are affine (in (x, u, w));
- the coupling functions θ_t^i are affine;
- the admissible set $\mathcal{U}_{i,t}^{\text{ad}} \neq \emptyset$ is a weak* closed, convex set;
- there exists an admissible control;
- the coupling constraint admits an optimal multiplier in L^2 .

For a step $\rho > 0$ small enough, the sequence of control generated by DADP converges in L^∞ toward the optimal control of the relaxed problem.

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Convergence result - Carpentier, L.

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- there exists an admissible control;
- the coupling constraint admits an optimal multiplier in L^1 .

For a step $\rho > 0$ small enough, there exists a subsequence of the sequence of control generated by DADP converging in L^∞ toward the optimal control of the relaxed problem.

Existence of Multiplier

Existence of multiplier -L.

Assume that

- the random noises \mathbf{W}_t are essentially bounded;
- the local cost functions L_t^i are finite and convex in (x_i, u_i) , continuous in w ;
- the dynamic functions f_t^i are affine in (x_i, u_i) , continuous in w ;
- the constraint functions θ_t^i are affine;
- there is no bound constraints on \mathbf{U}_t^i and \mathbf{X}_t^i .

Then, the coupling constraint admits a multiplier in \mathbb{L}^1 , hence the relaxed coupling constraint admits a multiplier in \mathbb{L}^1 .

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Bounds over the Original Problem

Upper and lower bounds

Lower Bound : For a given $\mu^{(k)}$ we have a lower bound of the dual of the relaxed problem $(\mathcal{P}^{\mathbf{Y}})$, hence a **lower bound** of the original problem (\mathcal{P}) .

Upper bound : Through an heuristic (using the DP equation) we can construct an admissible (for the original problem (\mathcal{P})) solution and hence obtain an upper bound (by Monte Carlo).

In practice, on a simple problem:

- around 3% gap with minimal information ($\mathbf{Y}_t \equiv 0$),
- around 2% gap with dynamic information.

Validity a posteriori

Validity

If we obtain a multiplier μ^\sharp leading to a solution $\mathbf{U}(\mu^\sharp)$ satisfying the (relaxed) constraint:

$$\mathbb{E} \left[\sum_{i=1}^N \theta_t(\mathbf{U}_t^i(\mu^\sharp)) \mid \mathbf{Y}_t \right] = 0$$

then the solution $\mathbf{U}(\mu^\sharp)$ is optimal (for the relaxed problem (\mathcal{P}_Y)).

Consequences:

- A Posteriori conclusion even if abstract conditions not verified,
- use of improved multiplier update step.

Conclusion of Part II

- Summing up DADP:
 - Choose an information process \mathbf{Y} following $\mathbf{Y}_{t+1} = \tilde{f}_t(\mathbf{Y}_t, \mathbf{W}_{t+1})$.
 - We relax the almost sure coupling constraint into a conditional expectation one and apply a price decomposition scheme to the relaxed problem.
 - The subproblems can be solved by dynamic programming with the state $(\mathbf{X}_t^i, \mathbf{Y}_t)$.
- We give:
 - a consistency result (family of information process),
 - a convergence result (fixed information process),
 - an existence of multiplier condition.

Thesis Outline

- 1 Preliminaries
- 2 Time-Consistency: from Optimization to Risk Measures
- 3 Stochastic Dual Dynamic Programming Algorithm
- 4 Constraint Qualification in Stochastic Optimization
- 5 Constraint Qualification in (L^∞, L^1)
- 6 Uzawa Algorithm in L^∞
- 7 Epiconvergence of Relaxed Stochastic Problems
- 8 Dual Approximate Dynamic Programming Algorithm

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Papers

-  P. Girardeau, V. Leclère and A. Philpott
On the convergence of decomposition methods for multi-stage stochastic convex programs.
Accepted in Mathematics of Operations Research, 2014
-  M. De Lara and V. Leclère
Time-Consistency: from Optimization to Risk Measures
Submitted, 2014
-  V. Leclère
Epiconvergence of relaxed stochastic optimization problem.
Submitted, 2013
-  M. Grasselli, M. Ludkovski and V. Leclère
Priority option: the value of being a leader.
IJTAF, 16, 2013

Conclusion: the next steps

2 Dynamic Programming

- extension of state
- more generic links

3 SDDP

- noise with compact support
- convergence estimation

4 L^1 multiplier

- bounds on control via Relatively Complete Recourse
- conditions for L^2 multiplier

6 Uzawa in L^∞

- reflexions around the strong-convexity
- use ε -convergence theory

7 Epiconvergence

- obtain a non-asymptotical bound

8 DADP

- Numerical test on big scale
- Method to construct Y
- Interactions with SDDP

The end

Thank you for your attention!