

Spatial Decomposition in Stochastic Optimization: Theoretical and Practical Questions.

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SESO Week

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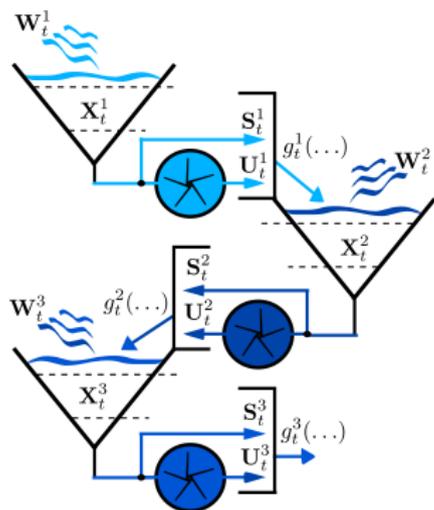
Multistage Stochastic Optimization : an Example

How to manage a chain of dam producing electricity from the turbine water to optimize the gain ?

$$\mathbb{E} \left[\sum_{i=1}^N \sum_{t=0}^{T-1} L_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}^i) + K_i(\mathbf{X}_T^i) \right]$$

Constraints:

- **dynamics:**
 $\mathbf{X}_{t+1}^i = f_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}^i),$
- **nonanticipativity:**
 $\mathbf{U}_t^i \preceq \mathcal{F}_t^i,$
- **spatial coupling:**
 $\mathbf{Z}_t^{i+1} = g_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}^i).$



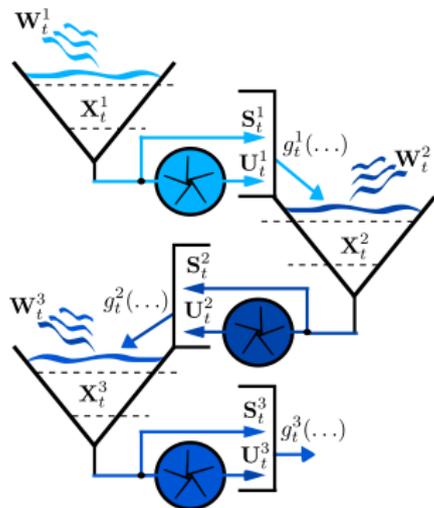
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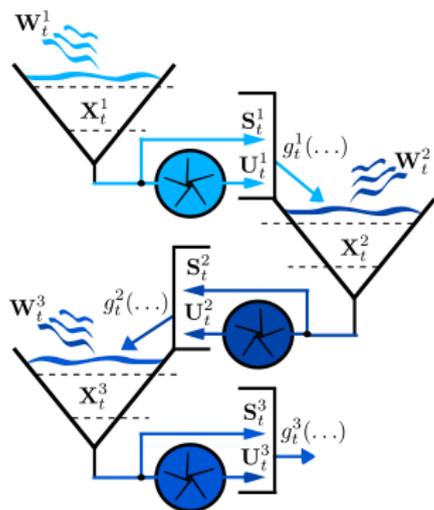
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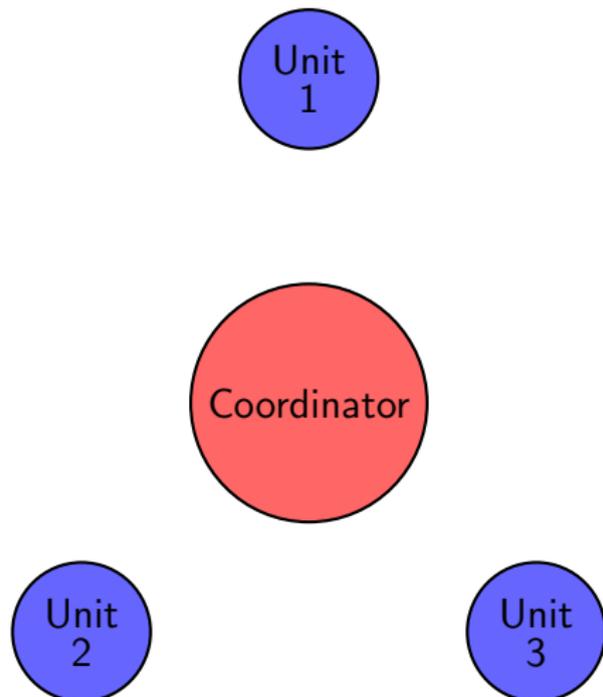
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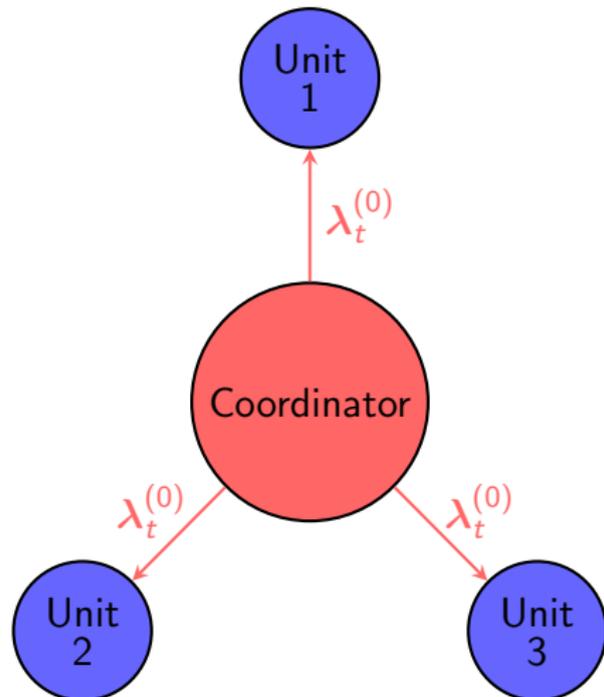
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- Satisfy a demand with N units of production at minimal cost.
- **Price decomposition:**
 - the coordinator sets a sequence of price λ_t ,
 - the units send their production planning $\mathbf{u}_t^{(i)}$,
 - the coordinator compare total production and demand and updates the price,
 - and so on...



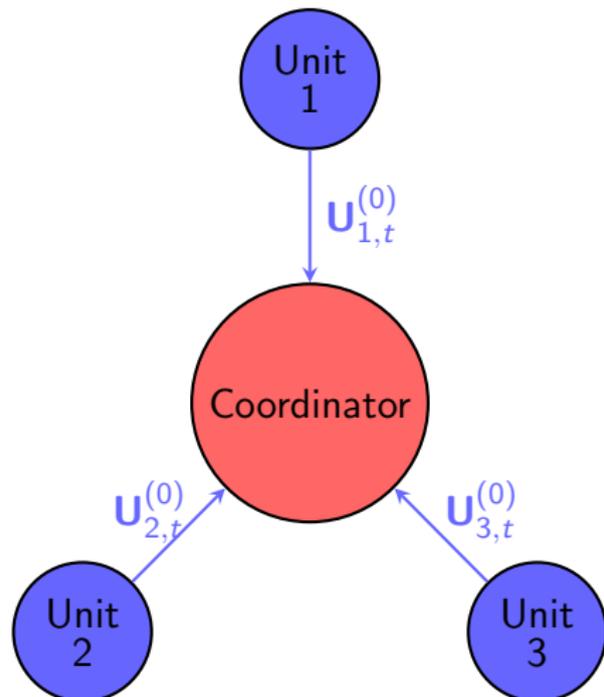
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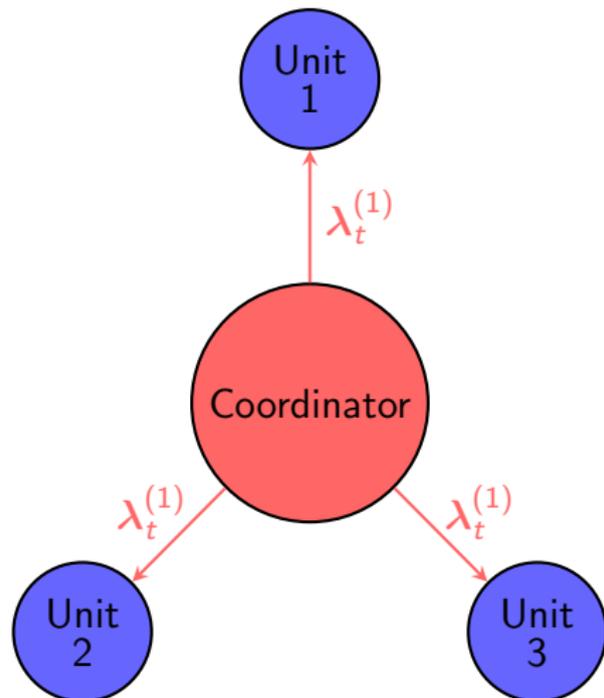
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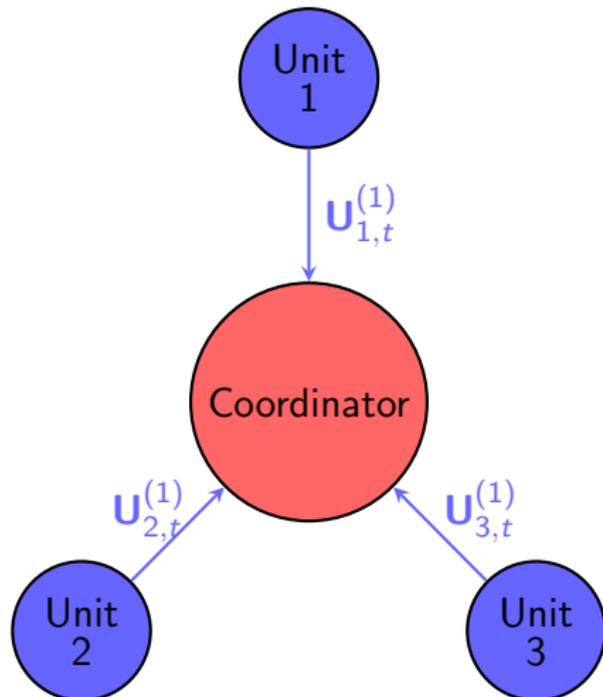
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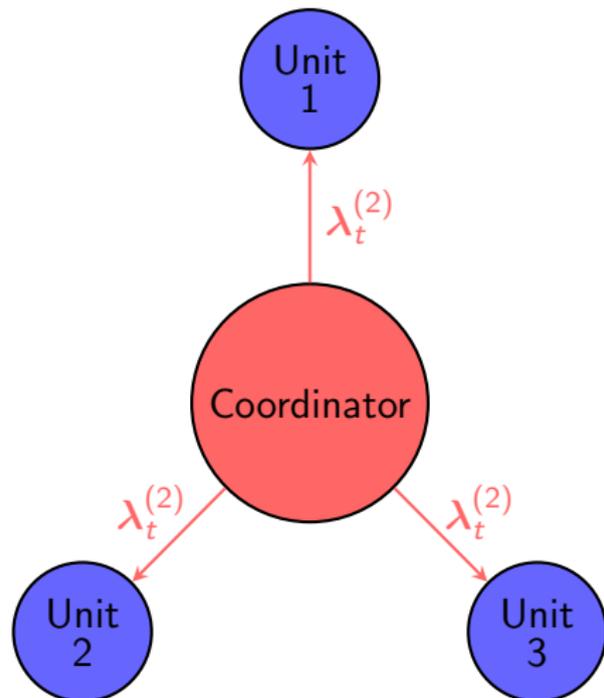
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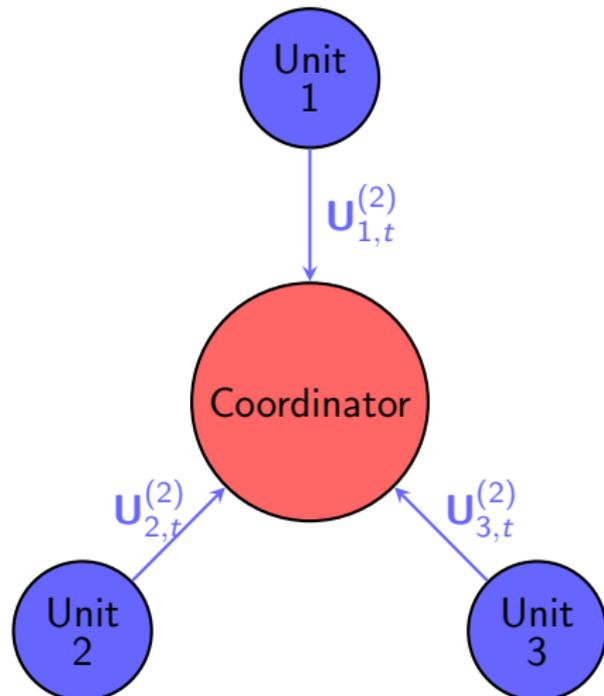
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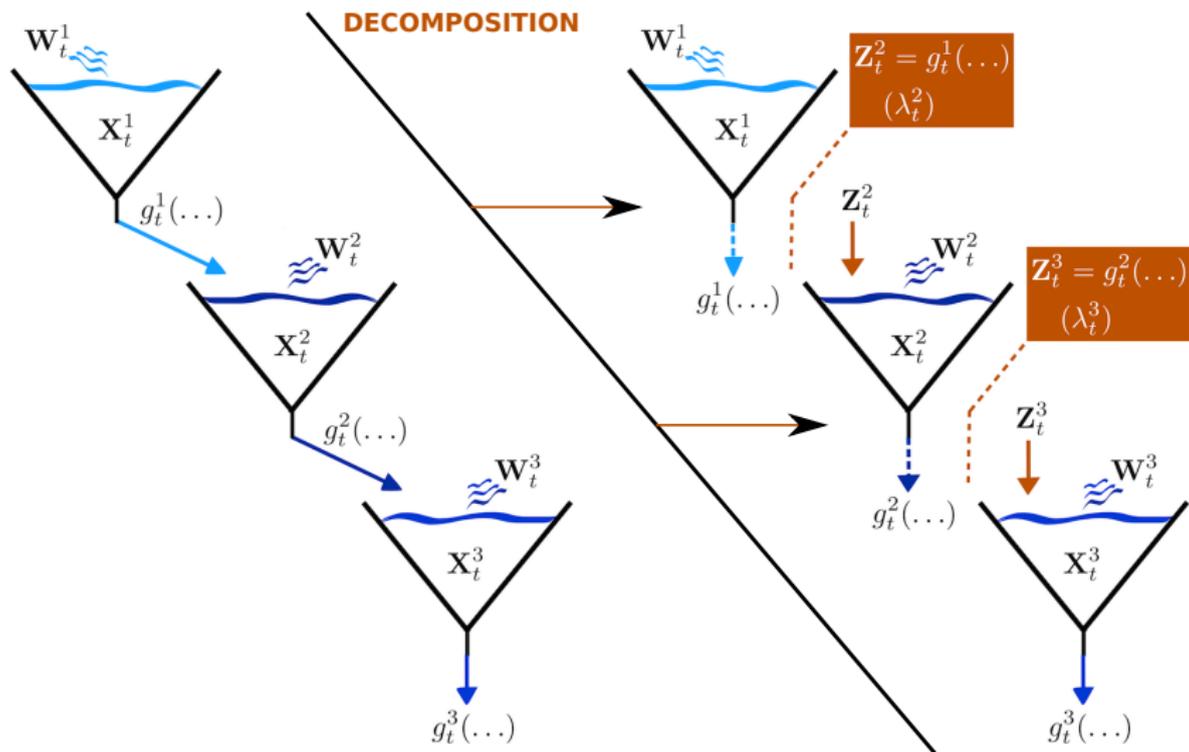
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Primal Problem

$$\begin{aligned} \min_{\mathbf{x}, \mathbf{u}} \quad & \sum_{i=1}^N \mathbb{E} \left[\sum_{t=0}^T L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + K^i(\mathbf{x}_T^i) \right] \\ \forall i, \quad & \mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}), \quad \mathbf{x}_0^i = \mathbf{x}_0^i, \\ \forall i, \quad & \mathbf{u}_t^i \in \mathcal{U}_{t,i}^{ad}, \quad \mathbf{u}_t^i \preceq \mathcal{F}_t, \\ & \sum_{i=1}^N \theta_t^i(\mathbf{u}_t^i) = 0 \end{aligned}$$

Solvable by DP with state $(\mathbf{x}_1, \dots, \mathbf{x}_N)$

Primal Problem

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 \forall i, \quad & \mathbf{u}_t^i \in \mathcal{U}_{t,i}^{ad}, \quad \mathbf{u}_t^i \preceq \mathcal{F}_t, \\
 \sum_{i=1}^N \theta_t^i(\mathbf{u}_t^i) = 0 \quad & \rightsquigarrow \boldsymbol{\lambda}_t \text{ multiplier}
 \end{aligned}$$

Solvable by DP with state $(\mathbf{x}_1, \dots, \mathbf{x}_N)$

Primal Problem with Dualized Constraint

$$\min_{\mathbf{x}, \mathbf{U}} \max_{\lambda} \sum_{i=1}^N \mathbb{E} \left[\sum_{t=0}^T L_t^i(\mathbf{x}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}) + \langle \lambda_t, \theta_t^i(\mathbf{U}_t^i) \rangle + K^i(\mathbf{x}_T^i) \right]$$

$$\forall i, \quad \mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}), \quad \mathbf{x}_0^i = \mathbf{x}_0^i,$$

$$\forall i, \quad \mathbf{U}_t^i \in \mathcal{U}_{t,i}^{ad}, \quad \mathbf{U}_t^i \preceq \mathcal{F}_t,$$

Coupling constraint dualized \implies all constraints are unit by unit

Dual Problem

$$\max_{\lambda} \min_{\mathbf{x}, \mathbf{u}} \sum_{i=1}^N \mathbb{E} \left[\sum_{t=0}^T L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + \langle \lambda_t, \theta_t^i(\mathbf{u}_t^i) \rangle + K^i(\mathbf{x}_T^i) \right]$$

$$\forall i, \quad \mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}), \quad \mathbf{x}_0^i = \mathbf{x}_0^i,$$

$$\forall i, \quad \mathbf{u}_t^i \in \mathcal{U}_{t,i}^{ad}, \quad \mathbf{u}_t^i \preceq \mathcal{F}_t,$$

Exchange operator **min** and **max** to obtain a new problem

Decomposed Dual Problem

$$\max_{\lambda} \sum_{i=1}^N \min_{\mathbf{x}^i, \mathbf{u}^i} \mathbb{E} \left[\sum_{t=0}^T L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + \langle \lambda_t, \theta_t^i(\mathbf{u}_t^i) \rangle + K^i(\mathbf{x}_T^i) \right]$$
$$\mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}), \quad \mathbf{x}_0^i = \mathbf{x}_0^i,$$
$$\mathbf{u}_t^i \in \mathcal{U}_{t,i}^{ad}, \quad \mathbf{u}_t^i \preceq \mathcal{F}_t,$$

For a given λ , minimum of sum is sum of minima

Inner Minimization Problem

$$\min_{\mathbf{x}^i, \mathbf{u}^i} \mathbb{E} \left[\sum_{t=0}^T L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + \langle \lambda_t, \theta_t^i(\mathbf{u}_t^i) \rangle + K^i(\mathbf{x}_T^i) \right]$$

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We have N smaller subproblems. Can they be solved by DP ?

Inner Minimization Problem

$$\min_{\mathbf{x}^i, \mathbf{u}^i} \mathbb{E} \left[\sum_{t=0}^T L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + \langle \boldsymbol{\lambda}_t, \theta_t^i(\mathbf{u}_t^i) \rangle + K^i(\mathbf{x}_T^i) \right]$$

$$\mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}), \quad \mathbf{x}_0^i = \mathbf{x}_0^i,$$

$$\mathbf{u}_t^i \in \mathcal{U}_{t,i}^{ad}, \quad \mathbf{u}_t^i \preceq \mathcal{F}_t,$$

No : $\boldsymbol{\lambda}$ is a time-dependent noise \rightsquigarrow state $(\mathbf{w}_1, \dots, \mathbf{w}_t)$

A Few Questions

- In which space lives the multiplier process λ ? For which duality ?
 - L^2
 - L^1
 - $(L^\infty)^*$
- What are the relations between the primal and dual problems ?
- Can we solve the subproblems by Dynamic Programming ?
 \rightsquigarrow No!
- How to update the multiplier process ?
 Uzawa Algorithm

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Problem Statement

We consider the following (primal) problem:

$$\begin{aligned}
 (\mathcal{P}) \quad & \min_{u \in \mathcal{U}^{\text{ad}}} J(u), \\
 & \text{s.t. } \Theta(u) \in -C.
 \end{aligned}$$

Where \mathcal{U} and \mathcal{V} are two Hausdorff spaces, and

- $J : \mathcal{U} \rightarrow \bar{\mathbb{R}}$ is an objective function ,
- $\Theta : \mathcal{U} \rightarrow \mathcal{V}$ is a constraint function (to be dualized),
- $C \subset \mathcal{V}$ is a cone of constraints,
- $\mathcal{U}^{\text{ad}} \subset \mathcal{U}$ is a constraint set (not to be dualized).

Dual Problem

The primal problem can be written

$$(\mathcal{P}) \quad \min_{u \in \mathcal{U}^{\text{ad}}} \max_{\lambda \in C^*} J(u) + \langle \lambda, \Theta(u) \rangle_{\mathcal{V}^*, \mathcal{V}},$$

where $C^* \subset \mathcal{V}^*$ is given by

$$C^* = \{ \lambda \in \mathcal{V}^* \mid \forall x \in C, \langle \lambda, x \rangle_{\mathcal{V}^*, \mathcal{V}} \geq 0 \}.$$

The dual problem of Problem (\mathcal{P}) reads

$$(\mathcal{D}) \quad \max_{\lambda \in C^*} \min_{u \in \mathcal{U}^{\text{ad}}} J(u) + \langle \lambda, \Theta(u) \rangle_{\mathcal{V}^*, \mathcal{V}}.$$

Gradient of the Dual

Assume that $\mathcal{U} = \mathcal{U}^*$, and $\mathcal{V} = \mathcal{V}^*$ are Hilbert spaces.
Recall the dual problem (\mathcal{D}) as

$$\max_{\lambda \in C^*} \underbrace{\min_{u \in \mathcal{U}^{\text{ad}}} \left\{ J(u) + \langle \lambda, \Theta(u) \rangle_{\mathcal{V}^*, \mathcal{V}} \right\}}_{:= \varphi(\lambda)} .$$

Under some regularity and unicity conditions, if $u^\sharp(\lambda)$ is a minimizer of the above problem, then

$$\Theta(u^\sharp(\lambda)) = \nabla \varphi(\lambda) .$$

Uzawa Algorithm

Data: Initial multiplier $\lambda^{(0)} \in \mathcal{V}$, step $\rho > 0$;

Result: Optimal solution $u^\#$ and multiplier $\lambda^\#$;

repeat

$$u^{(k)} \in \arg \min_{u \in \mathcal{U}^{\text{ad}}} \left\{ J(u) + \langle \lambda^{(k)}, \Theta(u) \rangle \right\},$$

$$\lambda^{(k+1)} = \text{proj}_{\mathcal{C}^*} \left(\lambda^{(k)} + \rho \Theta(u^{(k)}) \right).$$

until $\Theta(u^{(k)}) \in -\mathcal{C}$;

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L^∞ setting

From now on we consider that

$$\mathcal{U} = L^\infty(\Omega, \mathcal{F}, \mathbb{P}; \mathbb{R}^n) ,$$

$$\mathcal{V} = L^\infty(\Omega, \mathcal{F}, \mathbb{P}; \mathbb{R}^m) ,$$

$$\mathcal{C} = \{0\}.$$

Where the σ -algebra is not finite (modulo \mathbb{P}). Hence, \mathcal{U} and \mathcal{V} are non-reflexive, non-separable, Banach spaces.

If the σ -algebra is finite modulo \mathbb{P} , \mathcal{U} and \mathcal{V} are finite dimensional spaces, and the usual result applies.

Perks of an Hilbert Space

Fact

In an Hilbert space \mathcal{H} we know that

- i) the weak and weak* topologies are identical,
- ii) the space \mathcal{H} and its topological dual can be identified.

Point *i)* allows to formulate existence of minimizer results:

- weakly closed bounded \implies weakly compact;
- for a convex set: weakly closed \iff closed;
- for a convex function: weakly l.s.c \iff l.s.c.

Hence, a coercive, l.s.c. function J admits an infimum.

Point *ii)* allows to write gradient-like algorithm: at any iteration k , linear combination of $\lambda^{(k)}$ and $g^{(k)}$ take place in \mathcal{H} .

Difficulties Appearing in a Banach Space

- Reflexive Banach space:
 - *i*) still holds true (\rightsquigarrow existence of minimizers)
 - *ii*) no longer true (\rightsquigarrow linear combination of $u^{(k)} \in E$ and $g^{(k)} \in E^*$ does not have any sense).
- Non-reflexive Banach space E : neither *i*) nor *ii*) holds true.
- E is the topological dual of a Banach space: a weakly* closed bounded subset of E is weak* compact. Thus, weak* lower semicontinuity and coercivity of a function J gives the existence of minimizers of J .

Specificities of $L^\infty(\Omega, \mathcal{F}, \mathbb{P}; \mathbb{R}^n)$

- L^∞ is the topological dual of the Banach space L^1 . Hence, if J is weak* l.s.c and coercive, then J admits a minimizer.
- L^∞ can be identified with a subset of its topological dual $(L^\infty)^*$. Thus,

$$\lambda^{(k+1)} = \lambda^{(k)} + \rho \Theta(\mathbf{u}^{(k)}),$$

make sense: it is a linear combination of elements of $(L^\infty)^*$.

- Moreover, if $\lambda^{(0)}$ is chosen in L^∞ , then the sequence $\{\lambda^{(k)}\}_{k \in \mathbb{N}}$ remains in L^∞ .

Uzawa Algorithm

Data: Initial multiplier $\lambda^{(0)} \in L^\infty$, step $\rho > 0$;

Result: Optimal solution $\mathbf{U}^\#$ and multiplier $\lambda^\#$;

repeat

$$\begin{aligned} \mathbf{U}^{(k)} &\in \arg \min_{\mathbf{U} \in \mathcal{U}^{\text{ad}}} \left\{ J(\mathbf{U}) + \langle \lambda^{(k)}, \Theta(\mathbf{U}) \rangle \right\}, \\ \lambda^{(k+1)} &= \lambda^{(k)} + \rho \Theta(\mathbf{U}^{(k)}) . \end{aligned}$$

until $\Theta(\mathbf{U}^{(k)}) = 0$;

Remark: numerically, other update rules (e.g. quasi-Newton) can be used, convergence being proven when we find a multiplier $\lambda^{(k)}$ such that $\Theta(\mathbf{U}^{(k)}) = 0$.

Existence of Solution

Theorem

Assume that:

- 1 the constraint set \mathcal{U}^{ad} is weakly* closed;
- 2 $\Theta : \mathcal{U} \rightarrow \mathcal{V}$ is affine, weakly* continuous;
- 3 the objective function $J : \mathcal{U} \rightarrow \bar{\mathbb{R}}$ is weak* lower semicontinuous and coercive on \mathcal{U}^{ad} ;
- 4 there exists an admissible control.

Then the primal problem admits at least one solution.

Moreover for any $\lambda \in L^\infty(\Omega, \mathcal{F}, \mathbb{P}; \mathbb{R}^m)$

$$\arg \min_{\mathbf{U} \in \mathcal{U}^{\text{ad}}} \left\{ J(\mathbf{U}) + \langle \lambda, \Theta(\mathbf{U}) \rangle \right\} \neq \emptyset .$$

Convergence Result

Theorem

Assume that:

- ① $J : \mathcal{U} \rightarrow \bar{\mathbb{R}}$ is a proper, weak* lower semicontinuous, Gâteaux-differentiable, a -convex function;
- ② $\Theta : \mathcal{U} \rightarrow \mathcal{V}$ is affine, weak* continuous and κ -Lipschitz;
- ③ there exists an admissible control;
- ④ \mathcal{U}^{ad} is weak* closed convex;
- ⑤ there is an optimal L^1 -multiplier to the constraint $\Theta(\mathbf{U}) = 0$;
- ⑥ the step ρ is such that $0 < \rho < \frac{2a}{\kappa}$.

Then, Uzawa algorithm is well defined and there exists a subsequence $(\mathbf{U}^{(n_k)})_{k \in \mathbb{N}}$ converging in L^∞ toward the optimal solution $\mathbf{U}^\#$ of the primal problem.

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Standard duality in L^2 spaces

Assume that $\mathcal{U} = L^2(\Omega, \mathcal{A}, \mathbb{P}; \mathbb{R}^n)$ and $\mathcal{V} = L^2(\Omega, \mathcal{A}, \mathbb{P}; \mathbb{R}^m)$.

The standard sufficient **constraint qualification condition**

$$0 \in \text{ri}\left(\Theta(\mathcal{U}^{\text{ad}} \cap \text{dom}(J)) + C\right),$$

is **scarcely satisfied** in such a stochastic setting.

Proposition

If the σ -algebra \mathcal{A} is not finite modulo \mathbb{P} , then for any subset $U^{\text{ad}} \subset \mathbb{R}^n$ that is not an affine subspace, the set

$$U^{\text{ad}} = \left\{ \mathbf{U} \in L^p(\Omega, \mathcal{A}, \mathbb{P}; \mathbb{R}^n) \mid \mathbf{U} \in U^{\text{ad}} \quad \mathbb{P} - a.s. \right\},$$

has an empty relative interior in L^p , for any $p < +\infty$.

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has an empty relative interior in L^p , for any $p < +\infty$.

Standard duality in L^2 spaces

Consider the following optimization problem:

$$\begin{aligned} \inf_{u_0, \mathbf{U}_1} \quad & u_0^2 + \mathbb{E}[(\mathbf{U}_1 + \alpha)^2] , \\ \text{s.t.} \quad & u_0 \geq a , \\ & \mathbf{U}_1 \geq 0 , \\ & u_0 - \mathbf{U}_1 \geq \mathbf{W} , \end{aligned} \quad \text{to be dualized}$$

where \mathbf{W} is a random variable uniform on $[1, 2]$.

For $a < 2$:

- we can construct a maximizing sequence of multipliers for the dual problem that **does not converge** in L^2 ;
- this is a case of *non relatively complete recourse* (constraints on \mathbf{U}_1 induce stronger constraint on u_0);
- however there exists an optimal multiplier in $(L^\infty)^*$

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Constraint qualification in (L^∞, L^1)

From now on, we assume that

$$\begin{aligned} \mathcal{U} &= L^\infty(\Omega, \mathcal{A}, \mathbb{P}; \mathbb{R}^n) , \\ \mathcal{V} &= L^\infty(\Omega, \mathcal{A}, \mathbb{P}; \mathbb{R}^m) , \\ \mathcal{C} &= \{0\} , \end{aligned}$$

where the σ -algebra \mathcal{A} is not finite modulo \mathbb{P} .¹

We consider the **pairing** (L^∞, L^1) with the following topologies:

- $\sigma(L^\infty, L^1)$: weak* topology on L^∞ (**coarsest topology** such that all the L^1 -linear forms are continuous),
- $\tau(L^\infty, L^1)$: Mackey-topology on L^∞ (**finest topology** such that the continuous linear forms are only the L^1 -linear forms).

¹If the σ -algebra is finite modulo \mathbb{P} , \mathcal{U} and \mathcal{V} are finite dimensional spaces.

Weak* closedness of linear subspaces of L^∞

Proposition

Let $\Theta : L^\infty(\Omega, \mathcal{A}, \mathbb{P}; \mathbb{R}^n) \rightarrow L^\infty(\Omega, \mathcal{A}, \mathbb{P}; \mathbb{R}^m)$ be a linear operator, and assume that there exists a linear operator $\Theta^\dagger : L^1(\Omega, \mathcal{A}, \mathbb{P}; \mathbb{R}^m) \rightarrow L^1(\Omega, \mathcal{A}, \mathbb{P}; \mathbb{R}^n)$ such that:

$$\langle \mathbf{v}, \Theta(\mathbf{u}) \rangle = \langle \Theta^\dagger(\mathbf{v}), \mathbf{u} \rangle \quad \forall \mathbf{u}, \forall \mathbf{v} .$$

Then the linear operator Θ is weak* continuous.

Applications

- $\Theta(\mathbf{u}) = \mathbf{u} - \mathbb{E}[\mathbf{u} \mid \mathcal{B}]$: non-anticipativity constraints,
- $\Theta(\mathbf{u}) = A\mathbf{u}$ with $A \in \mathcal{M}_{m,n}(\mathbb{R})$: finite number of constraints.

A duality theorem

$$(\mathcal{P}) \quad \min_{\mathbf{U} \in \mathcal{U}} J(\mathbf{U}) \quad \text{s.t.} \quad \Theta(\mathbf{U}) = 0,$$

with $J(\mathbf{U}) = \mathbb{E}[j(\mathbf{U}, \mathbf{W})]$.

Theorem

Assume that j is a convex normal integrand, that J is continuous in the Mackey topology at some point \mathbf{U}_0 such that $\Theta(\mathbf{U}_0) = 0$, and that Θ is weak* continuous on $L^\infty(\Omega, \mathcal{A}, \mathbb{P}; \mathbb{R}^n)$.

Then, $\mathbf{U}^\# \in \mathcal{U}$ is an optimal solution of Problem (\mathcal{P}) if and only if there exists $\lambda^\# \in L^1(\Omega, \mathcal{A}, \mathbb{P}; \mathbb{R}^m)$ such that

- $\mathbf{U}^\# \in \arg \min_{\mathbf{U} \in \mathcal{U}} \mathbb{E} [j(\mathbf{U}, \mathbf{W}) + \lambda^\# \cdot \Theta(\mathbf{U})],$
- $\Theta(\mathbf{U}^\#) = 0.$

Extension of a result given by R. Wets for non-anticipativity constraints.

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Dual approximation as constraint relaxation

The original problem is (abstract form)

$$\begin{aligned} \min_{\mathbf{U} \in \mathcal{U}} \quad & J(\mathbf{U}) \\ \text{s.t.} \quad & \Theta(\mathbf{U}) = 0 \end{aligned}$$

written as

$$\min_{\mathbf{U} \in \mathcal{U}} \max_{\lambda} J(\mathbf{U}) + \mathbb{E}[\langle \lambda, \Theta(\mathbf{U}) \rangle]$$

Substituting λ by $\mathbb{E}(\lambda | \mathbf{Y})$ gives

$$\min_{\mathbf{U} \in \mathcal{U}} \max_{\lambda} J(\mathbf{U}) + \mathbb{E}[\langle \mathbb{E}(\lambda | \mathbf{Y}), \Theta(\mathbf{U}) \rangle]$$

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equivalent to

$$\begin{aligned} \min_{\mathbf{U} \in \mathcal{U}} \quad & J(\mathbf{U}) \\ \text{s.t.} \quad & \mathbb{E}(\Theta(\mathbf{U}) | \mathbf{Y}) = 0 \end{aligned}$$

Recall of the Multistage Problem

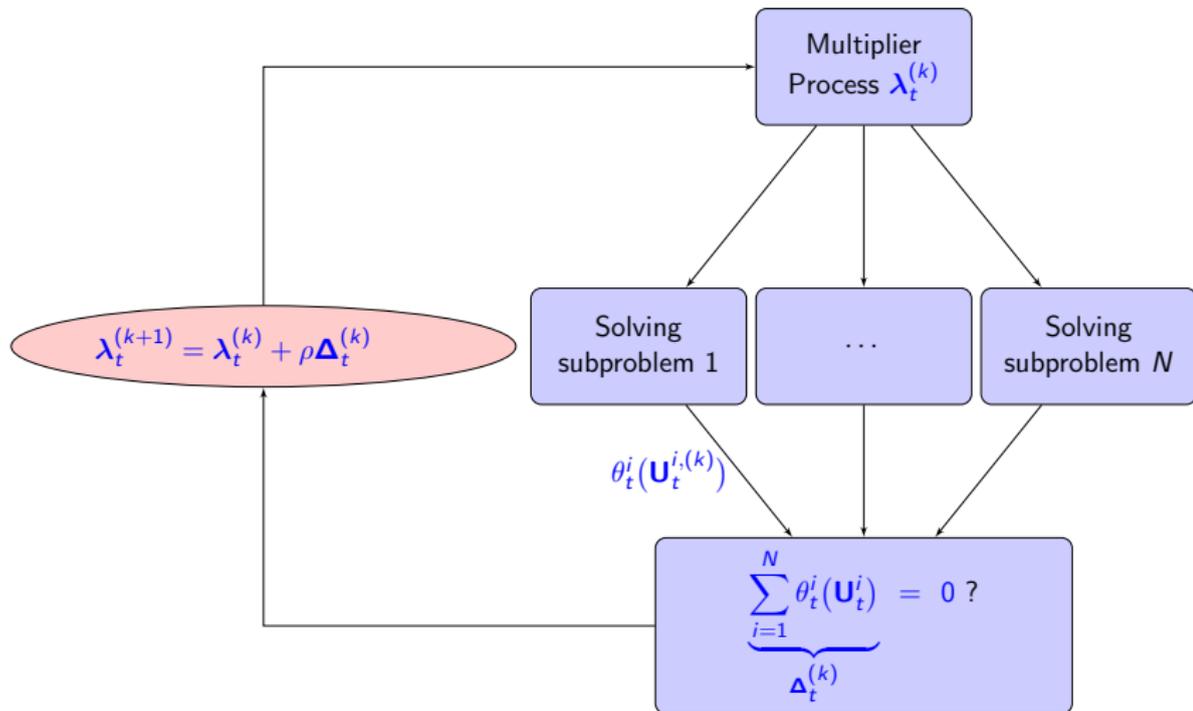
$$\min_{\mathbf{U}} \sum_{i=1}^N \mathbb{E} \left[\sum_{t=1}^{T-1} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + K^i(\mathbf{x}_T) \right]$$

$$\mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}), \quad \mathbf{x}_0^i = \mathbf{x}_0^i$$

$$\mathbf{u}_t^i \in \mathcal{U}_{t,i}^{ad}, \quad \mathbf{u}_t^i \preceq \mathcal{F}_t$$

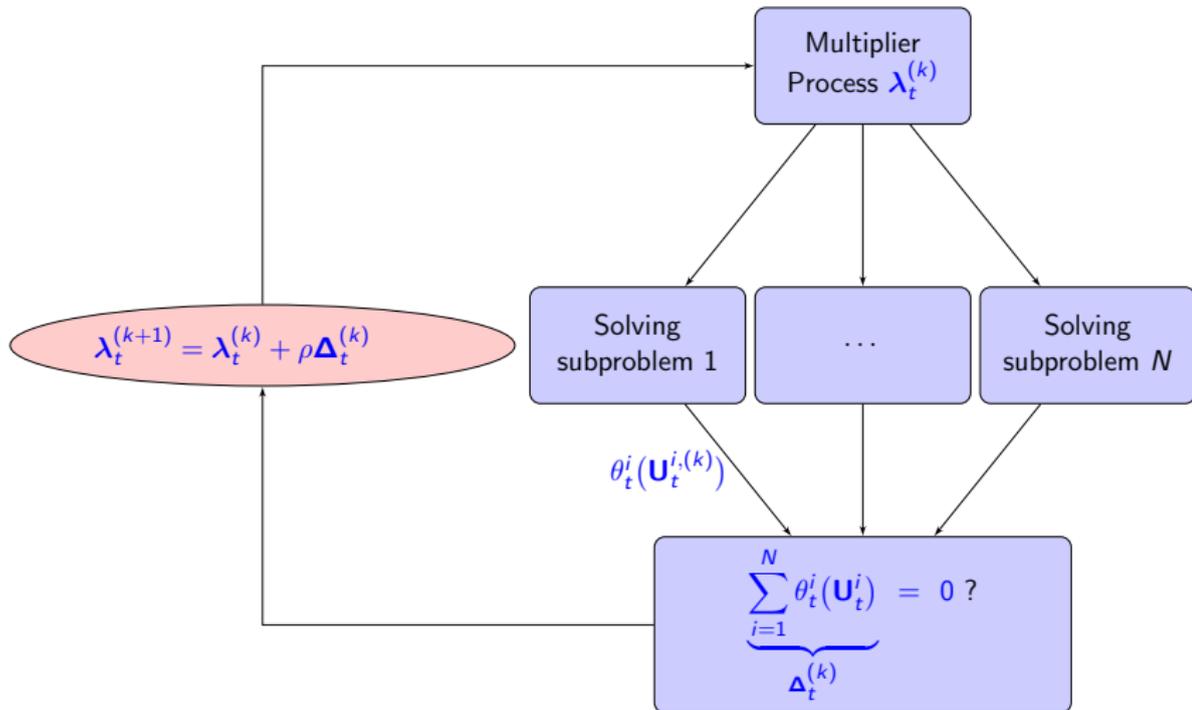
$$\sum_{i=1}^N \theta_t^i(\mathbf{u}_t^i) = 0 \quad \rightsquigarrow \boldsymbol{\lambda}_t$$

Stochastic spatial decomposition scheme



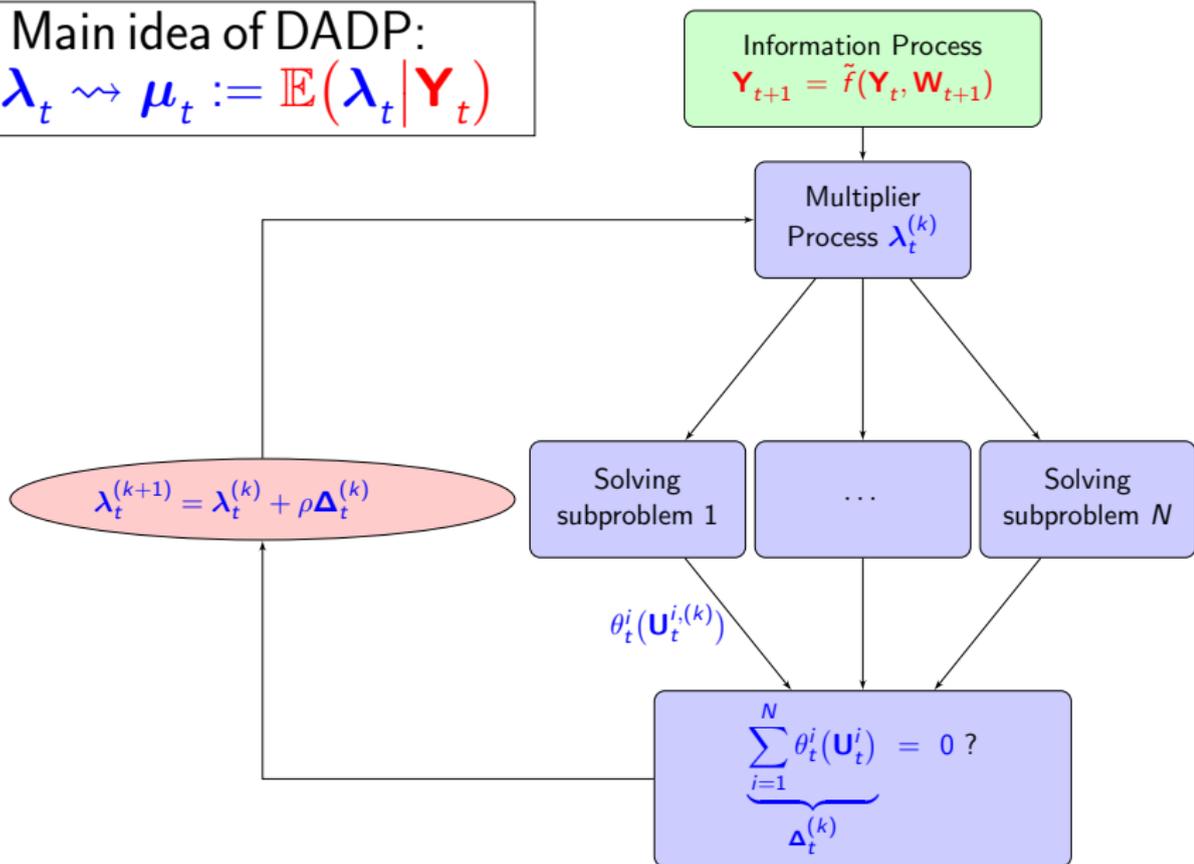
Main idea of DADP:

$$\lambda_t \rightsquigarrow \mu_t := \mathbb{E}(\lambda_t | \mathbf{Y}_t)$$

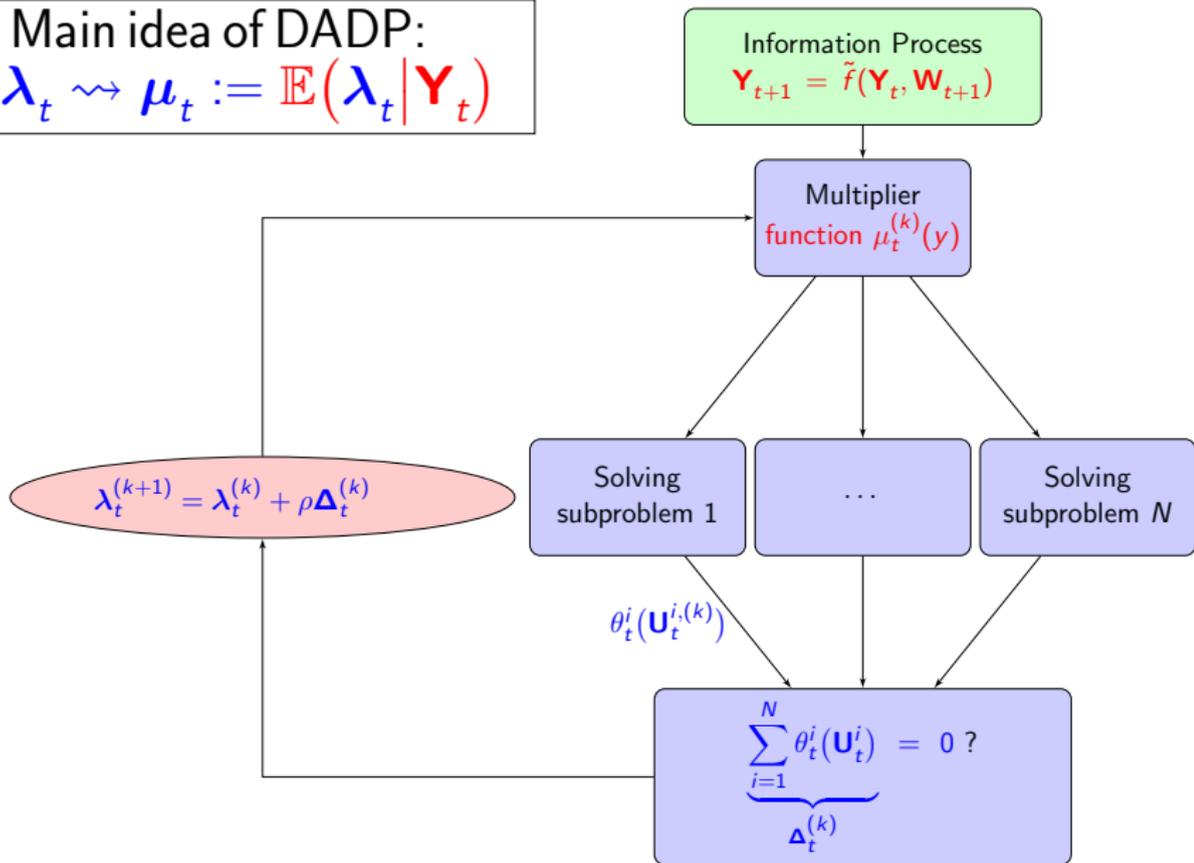


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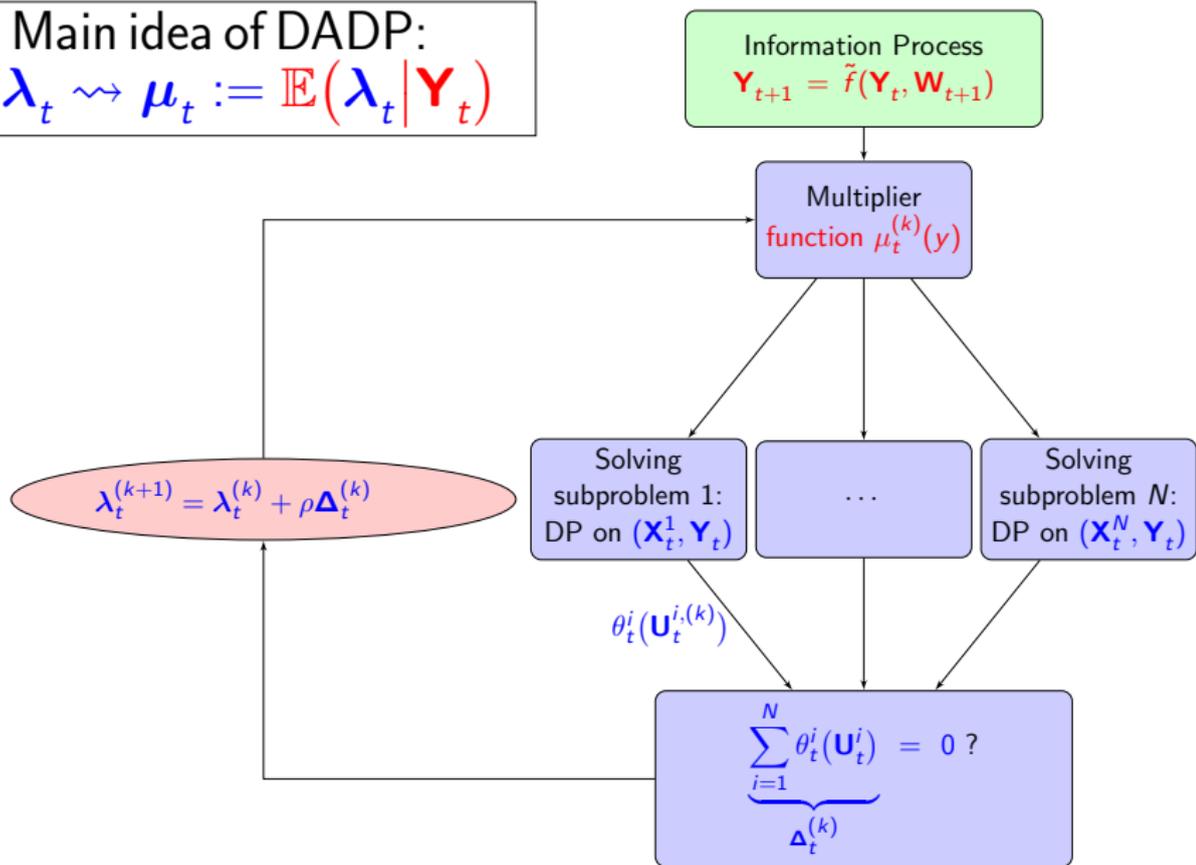
$$\lambda_t \rightsquigarrow \mu_t := \mathbb{E}(\lambda_t | \mathbf{Y}_t)$$



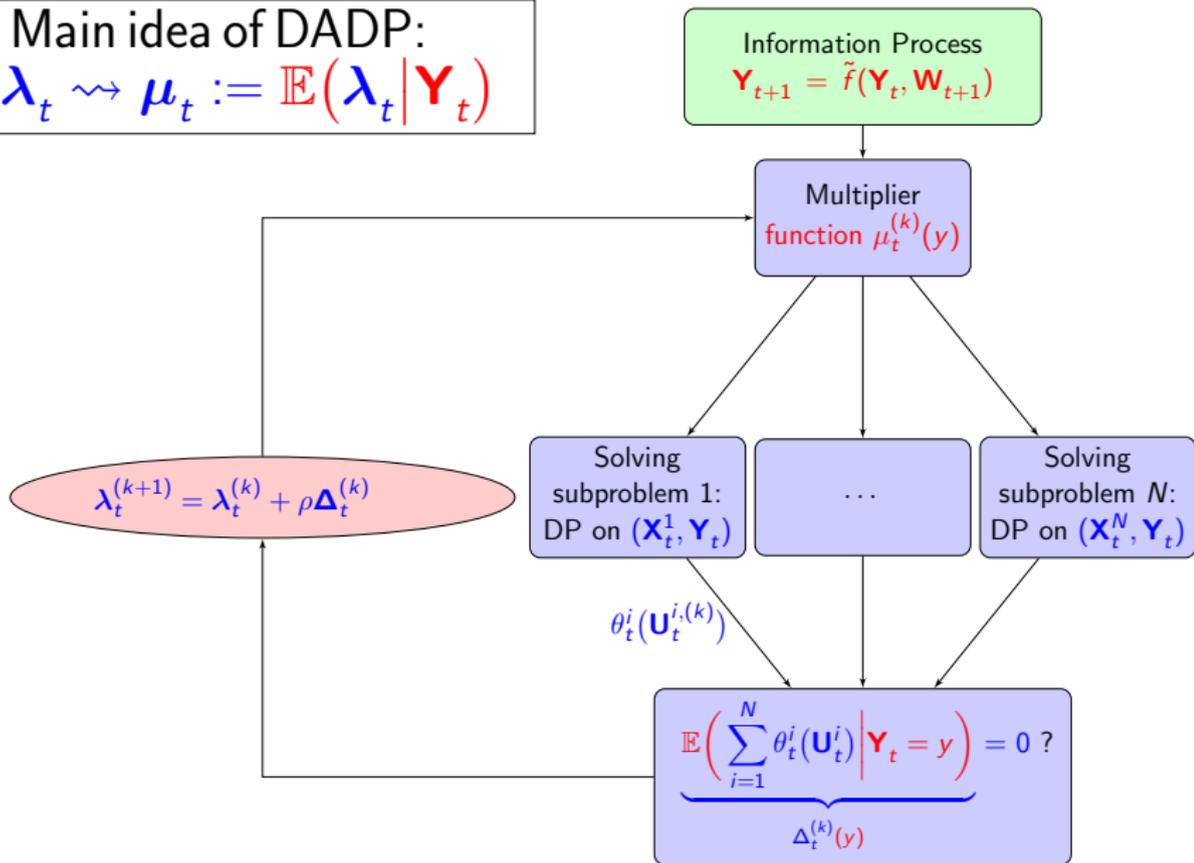
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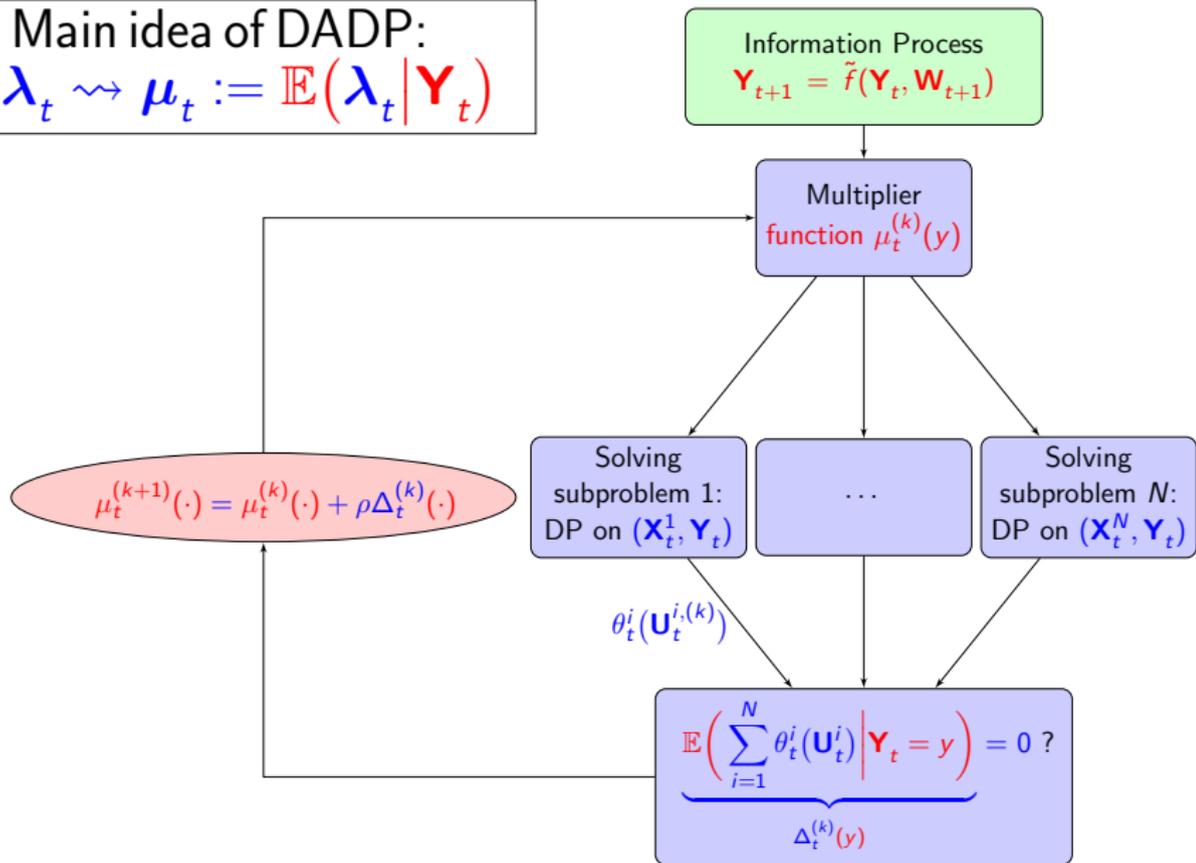
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Global problem:

$$\min_{\{\mathbf{u}_t^i\}_{i,t}} \sum_{i=1}^N \mathbb{E} \left[\sum_{t=1}^T L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + K(\mathbf{x}_T^i) \right]$$

$$\mathbf{x}_{t+1}^i = f_t(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}), \quad \mathbf{x}_0^i = \mathbf{x}_0^i$$

$$\mathbf{u}_t^i \in \mathcal{U}_{t,i}^{ad}, \quad \mathbf{u}_t^i \preceq \mathcal{F}_t$$

$$\sum_{i=1}^n \theta_t^i(\mathbf{u}_t^i) = 0$$

Solved by DP with state $(\mathbf{x}_t^1, \dots, \mathbf{x}_t^N)$:

$$V_t(\mathbf{x}) = \min_{\{\mathbf{u}_t^i\}_i} \sum_{i=1}^N \mathbb{E} \left[L_t^i(x_t^i, u_t^i, \mathbf{w}_{t+1}) + V_{t+1}(\mathbf{x}_{t+1}) \right]$$

$$\mathbf{x}_{t+1}^i = f_t(x_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}),$$

$$u_t^i \in \mathcal{U}_{t,i}^{ad},$$

$$\sum_{i=1}^n \theta_t^i(\mathbf{u}_t^i) = 0$$

Subproblem of Stochastic Decomposition

$$\min_{\{\mathbf{u}_t^i\}_t} \mathbb{E} \left[\sum_{t=1}^T L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + \langle \lambda_t, \theta_t(\mathbf{u}_t^i) \rangle + K(\mathbf{x}_T^i) \right]$$

$$\mathbf{x}_{t+1}^i = f_t(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}), \quad \mathbf{x}_0^i = \mathbf{x}_0^i$$

$$\mathbf{u}_t^i \in \mathcal{U}_{t,i}^{ad}, \quad \mathbf{u}_t^i \preceq \mathcal{F}_t$$

Solved by DP with state $(\mathbf{w}_1, \dots, \mathbf{w}_t)$:

$$V_t(\{\mathbf{w}_\tau\}_1^{t-1}) = \min_{\{\mathbf{u}_t^i\}} \mathbb{E} \left[L_t^i(x_t^i, u_t^i, \mathbf{w}_{t+1}) + \langle \lambda_t, \theta_t(\mathbf{u}_t^i) \rangle + V_{t+1}(\{\mathbf{w}_\tau\}_1^t) \right]$$

$$\{\mathbf{w}_\tau\}_1^{t-1} = \{\mathbf{w}_\tau\}_1^{t-1}$$

$$\mathbf{x}_{t+1}^i = f_t(x_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}),$$

$$u_t^i \in \mathcal{U}_{t,i}^{ad},$$

Subproblem of DADP

$$\min_{\{\mathbf{U}_t^i\}_t} \mathbb{E} \left[\sum_{t=1}^T L_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}) + \langle \mu_t(\mathbf{Y}_t), \theta_t(\mathbf{U}_t^i) \rangle + K(\mathbf{X}_T^i) \right]$$

$$\mathbf{X}_{t+1}^i = f_t(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}), \quad \mathbf{X}_0^i = \mathbf{x}_0^i$$

$$\mathbf{U}_t^i \in \mathcal{U}_{t,i}^{ad}, \quad \mathbf{U}_t^i \preceq \mathcal{F}_t$$

$$\mathbf{Y}_{t+1} = \tilde{f}_t(\mathbf{Y}_t, \mathbf{W}_{t+1})$$

Solved by DP with state $(\mathbf{X}_t^i, \mathbf{Y}_t)$:

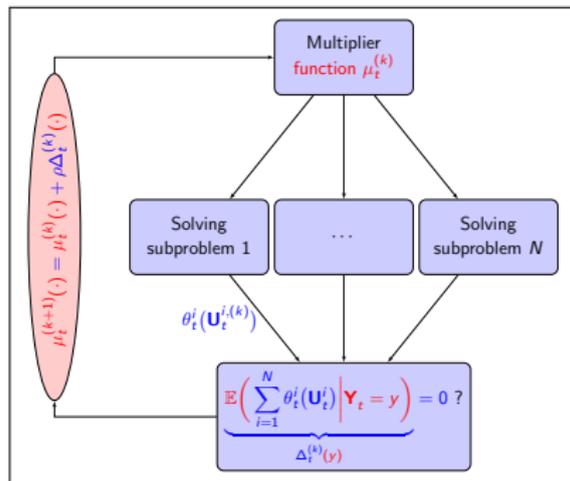
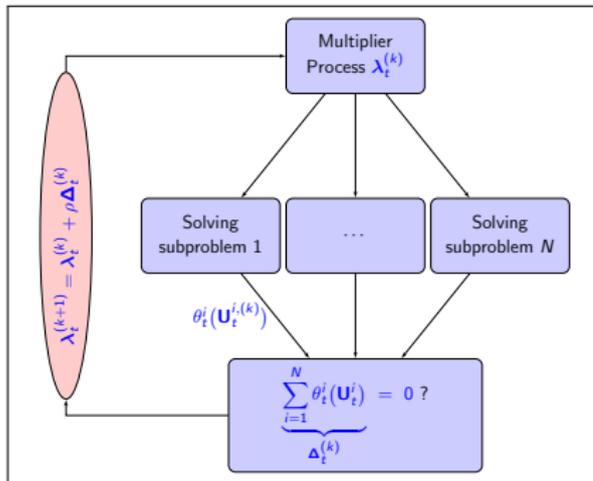
$$V_t^i(x, y) = \min_{\{\mathbf{U}_t^i\}} \mathbb{E} \left[L_t^i(x_t^i, u_t^i, \mathbf{W}_{t+1}) + \langle \mu_t(\mathbf{Y}_t), \theta_t(\mathbf{U}_t^i) \rangle + V_{t+1}(\mathbf{X}_{t+1}^i, \mathbf{Y}_{t+1}) \right]$$

$$\mathbf{X}_{t+1}^i = f_t(x_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}),$$

$$u_t^i \in \mathcal{U}_{t,i}^{ad},$$

$$\mathbf{Y}_{t+1} = \tilde{f}_t(y, \mathbf{W}_{t+1})$$

Main idea of DADP: $\lambda_t \rightsquigarrow \mu_t := \mathbb{E}(\lambda_t | \mathbf{Y}_t)$



Main problems:

- Subproblems not easily solvable by DP
- $\lambda^{(k)}$ live in a huge space

Advantages:

- Subproblems solvable by DP with state $(\mathbf{X}_t^i, \mathbf{Y}_t)$
- $\mu^{(k)}$ live in a smaller space

3 Interpretations of DADP

- DADP as an approximation of the optimal multiplier

$$\lambda_t \rightsquigarrow \mathbb{E}(\lambda_t | \mathbf{Y}_t) .$$

- DADP as a decision-rule approach in the dual

$$\max_{\lambda} \min_{\mathbf{U}} L(\lambda, \mathbf{U}) \rightsquigarrow \max_{\lambda_t \preceq \mathbf{Y}_t} \min_{\mathbf{U}} L(\lambda, \mathbf{U}) .$$

- DADP as a constraint relaxation

$$\sum_{i=1}^n \theta_t^i(\mathbf{u}_t^i) = 0 \rightsquigarrow \mathbb{E} \left(\sum_{i=1}^n \theta_t^i(\mathbf{u}_t^i) \middle| \mathbf{Y}_t \right) = 0 .$$

Theoretical Results

- Consistence of the approximation (if we consider a sequence of approximated problems).
- Existence of multiplier of the coupling constraint.
- Convergence of the decomposition algorithm for a given relaxation.
- Lower and upper bounds on the original problem.
- A posteriori verification allowing for better multiplier update.

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Choosing an Information Process \mathbf{Y}

- **Perfect memory:** $\mathbf{Y}_t^i = (\mathbf{W}_0, \dots, \mathbf{W}_t)$.
 \rightsquigarrow equivalent to original problem, no numerical gain.
- **Minimal information:** $\mathbf{Y}_t^i \equiv \text{cste}$.
 \rightsquigarrow equivalent to replacing a.s. constraint by expected constraint. Subproblems solved efficiently (state \mathbf{X}_t^i), multiplier is deterministic.
- **Static information:** $\mathbf{Y}_t^i = h_t^i(\mathbf{W}_t)$.
 \rightsquigarrow Subproblems solved efficiently (state \mathbf{X}_t^i).
- **Dynamic information:** $\mathbf{Y}_{t+1}^i = h_t^i(\mathbf{Y}_t^i, \mathbf{W}_{t+1})$.
 \rightsquigarrow A number of possibilities. Some ideas:
 - mimicking the trajectory of the state of another unit (phantom state),
 - mimicking the control of other units,
 - Markov chain representing roughly the general state of the system.

Numerical Advantages of a finitely supported \mathbf{Y}

- Assume that each noise \mathbf{W}_t take w values, and the constraint function take value in \mathbb{R} .
- Then the multiplier λ_t of the almost sure constraint at time t lives in \mathbb{R}^{wt} .
- Assume that the information process at time t take y values, then the multiplier of the relaxed constraint μ_t lives in \mathbb{R}^y .
- Moreover each subproblems take “only” roughly y times more computational effort to solve than the subproblem with local state \mathbf{X}_t^i .

Back to Admissibility

- Consider an information process $\mathbf{Y}_{t+1} = \tilde{f}_t(\mathbf{Y}_t, \mathbf{W}_{t+1})$.
- For a multiplier process $\mu_t^{(k)}$ we obtain local Bellman function

$$\begin{aligned} \tilde{V}_t^i(x^i, y) &= \min_{u^i} \mathbb{E} \left[L_t^i(x^i, u^i, \mathbf{W}_{t+1}) + \tilde{V}_t^i(x_{t+1}^i, y_{t+1}) \right] \\ \mathbf{X}_{t+1}^i &= f_t(x^i, u^i, \mathbf{W}_{t+1}) \\ \mathbf{Y}_{t+1} &= \tilde{f}_t(y, \mathbf{W}_{t+1}) \end{aligned}$$

- An admissible strategy is given by

$$\begin{aligned} \pi_t^{\text{ad}}(x, y) \in \arg \min_{\{u^i\}_{i \in [1, N]}} \mathbb{E} \left[\sum_{i=1}^N \left(L_t^i(x^i, u^i, \mathbf{W}_{t+1}) + \tilde{V}_t^i(\mathbf{X}_{t+1}, \mathbf{Y}_{t+1}) \right) \right] \\ \mathbf{X}_{t+1}^i = f_t^i(x^i, u^i, \mathbf{W}_{t+1}), \quad \forall i \\ \mathbf{Y}_{t+1} = \tilde{f}_t(y, \mathbf{W}_{t+1}) \end{aligned}$$

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Problem Specification

- We consider a 3 dam problem, over 12 time steps.
- We relax each constraint with a given information process \mathbf{Y}^i .
- All random variable are discrete (noise, control, state).
- We test the following information processes:
 - Constant information** equivalent to replace the a.s. constraint by an expected constraint,
 - Part of noise** the information process is the inflow of the above dam $\mathbf{Y}_t^i = \mathbf{W}_t^{i-1}$,
 - Phantom state** the information process mimick the optimal trajectory of the state of the first dam (by statistical regression over the known optimal trajectory in this case)

Numerical Results on the 3 Dams Example

	DADP - \mathbb{E}	DADP - \mathbf{W}^{i-1}	DADP - dyn.	DP
Nb of it.	165	170	25	1
Time (min)	2	3	67	41
Lower Bound	-1.386×10^6	-1.379×10^6	-1.373×10^6	
Final Value	-1.335×10^6	-1.321×10^6	-1.344×10^6	-1.366×10^6
Loss	-2.3%	-3.3%	-1.6%	ref.

Table: Numerical results on the 3-dam problem

Summing up DADP

- Choose an information process \mathbf{Y} following
$$\mathbf{Y}_{t+1} = \tilde{f}_t(\mathbf{Y}_t, \mathbf{W}_{t+1}).$$
- We relax the almost sure coupling constraint into a conditional expectation one and apply a price decomposition scheme to the relaxed problem.
- The subproblems can be solved by dynamic programming with the state $(\mathbf{X}_t^i, \mathbf{Y}_t)$.
- We give a consistency result (family of information process), a convergence result (fixed information process) and an existence of multiplier condition.

The end

Thanks for your attention!

More information² on theoretical results tomorrow at ENPC,
amphi Caquot I, (14h).

²and hopefully some champagne